



# Life in the Fast Lane .... With a Model-T

## Why Last Mile Capacity is Just the Beginning



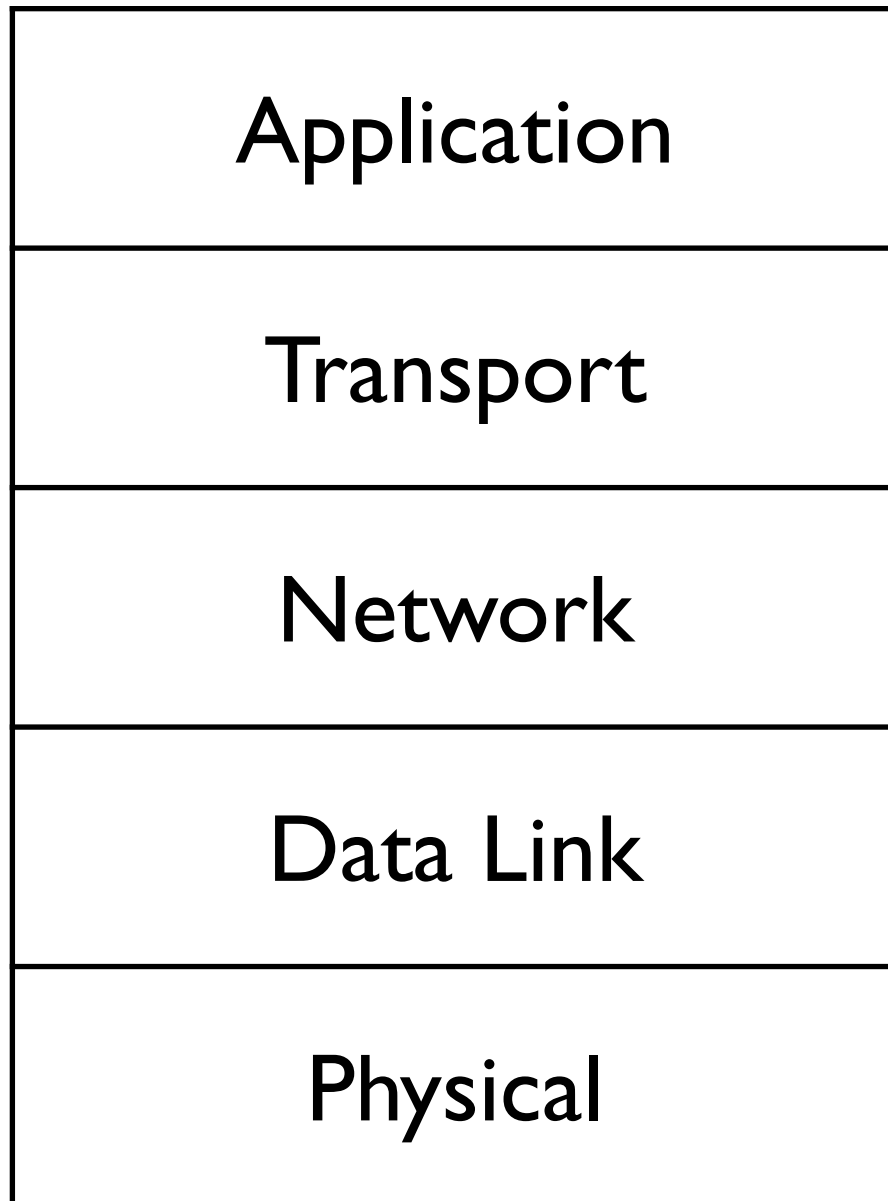
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*ICSI / Case*

Cyberinfrastructure Day  
Case Western Reserve University  
March 27, 2015

*“With the radio blatin’,  
Goin’ cruisin’ just as fast as we can now”*

# Thinking About Capacity

## Internet Layers



- This is how we theoretically and practically organize networks

# Thinking About Capacity

## Internet Layers

- This is what we consider when we discuss “capacity”
- Of course, layer 1 capacity is required



Physical

# Thinking About Capacity

## Internet Layers



Application

- But, this is where capacity really matters



# Thinking About Capacity

## Internet Layers

Application
Transport
Network
Data Link
Physical



# Wireless Networks

## Internet Layers

Application
Transport
Network
Data Link
Physical



# Residential Networks

## Internet Layers

Application
Transport
Network
Data Link
Physical



# Wired Enterprise Networks

## Internet Layers

Application
Transport
Network
Data Link
Physical



# Data Center Networks

## Internet Layers

Application
Transport
Network
Data Link
Physical



# Highly Specialized Networks

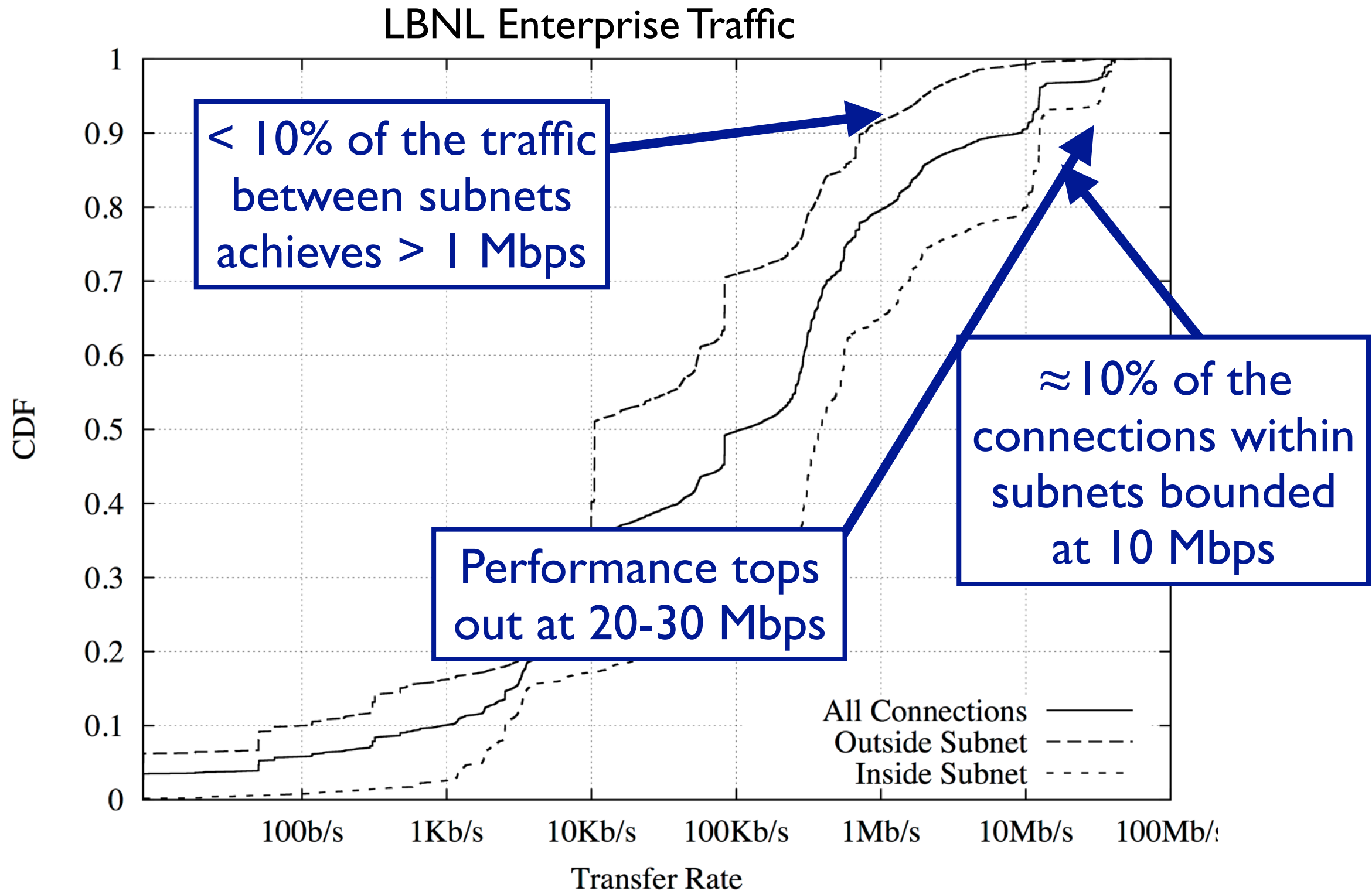
## Internet Layers

Application
Transport
Network
Data Link
Physical



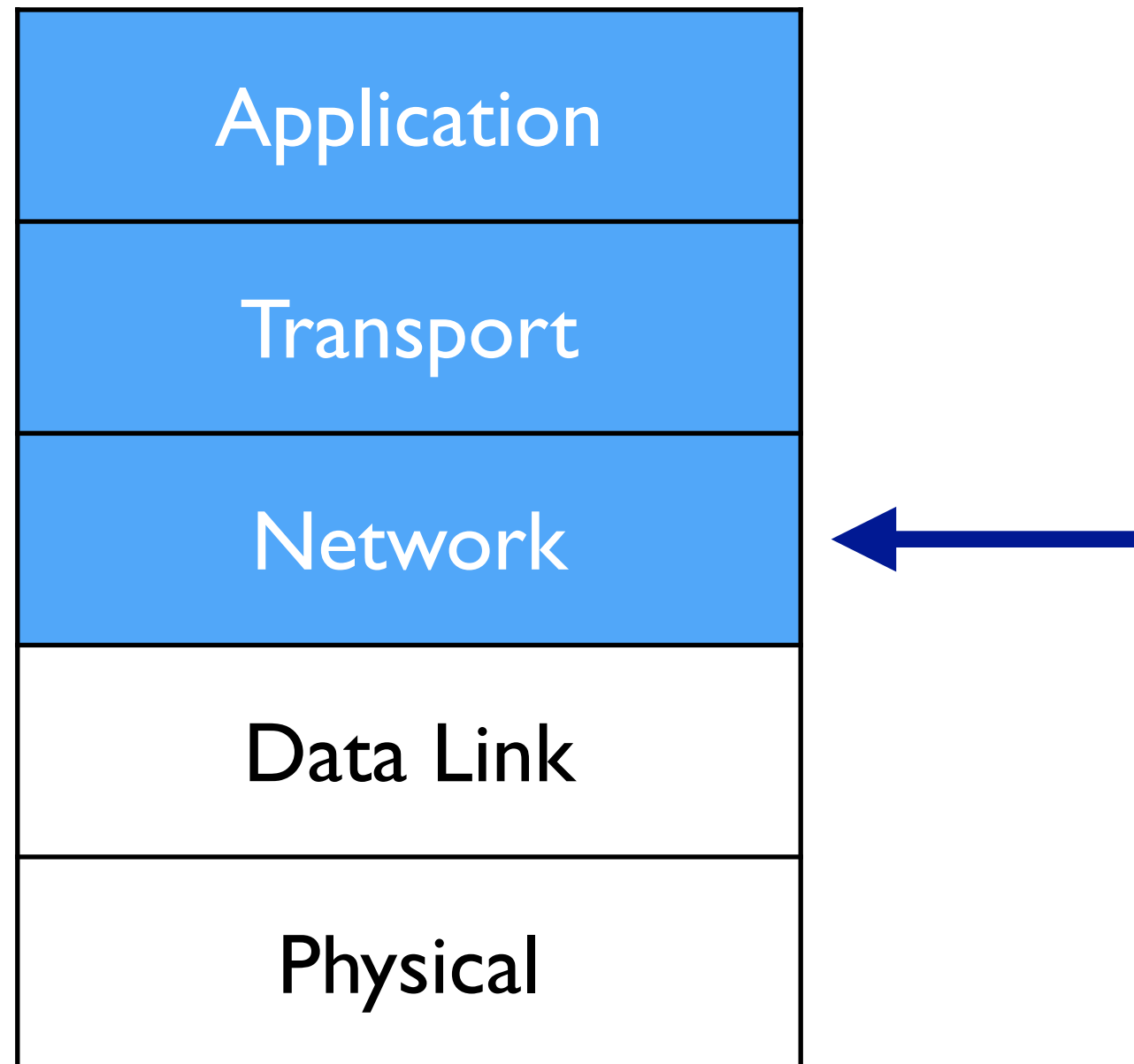
... but ...

# A First Look



# It's A Long Way To The Top ...

## Internet Layers

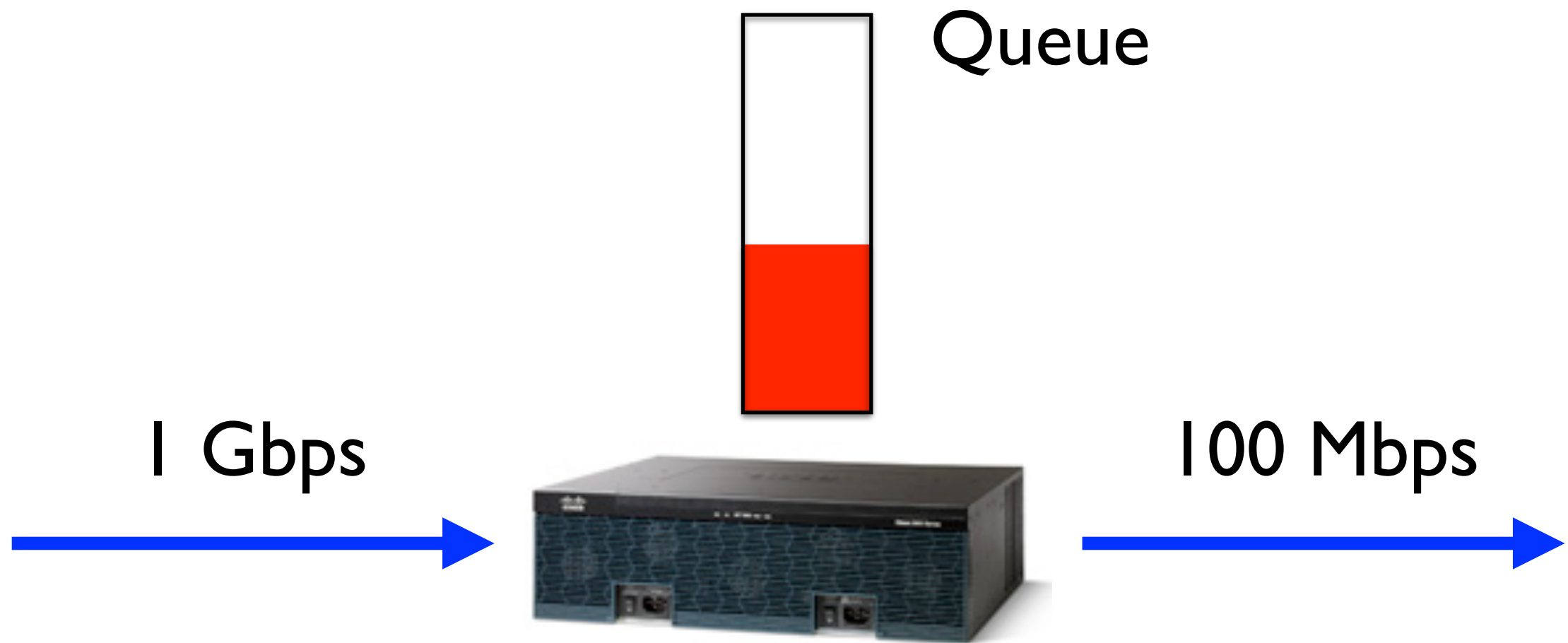




# Network Layer

- Fundamentally about moving packets from one host to another
- But, also, naturally must deal with *impedance mismatches*

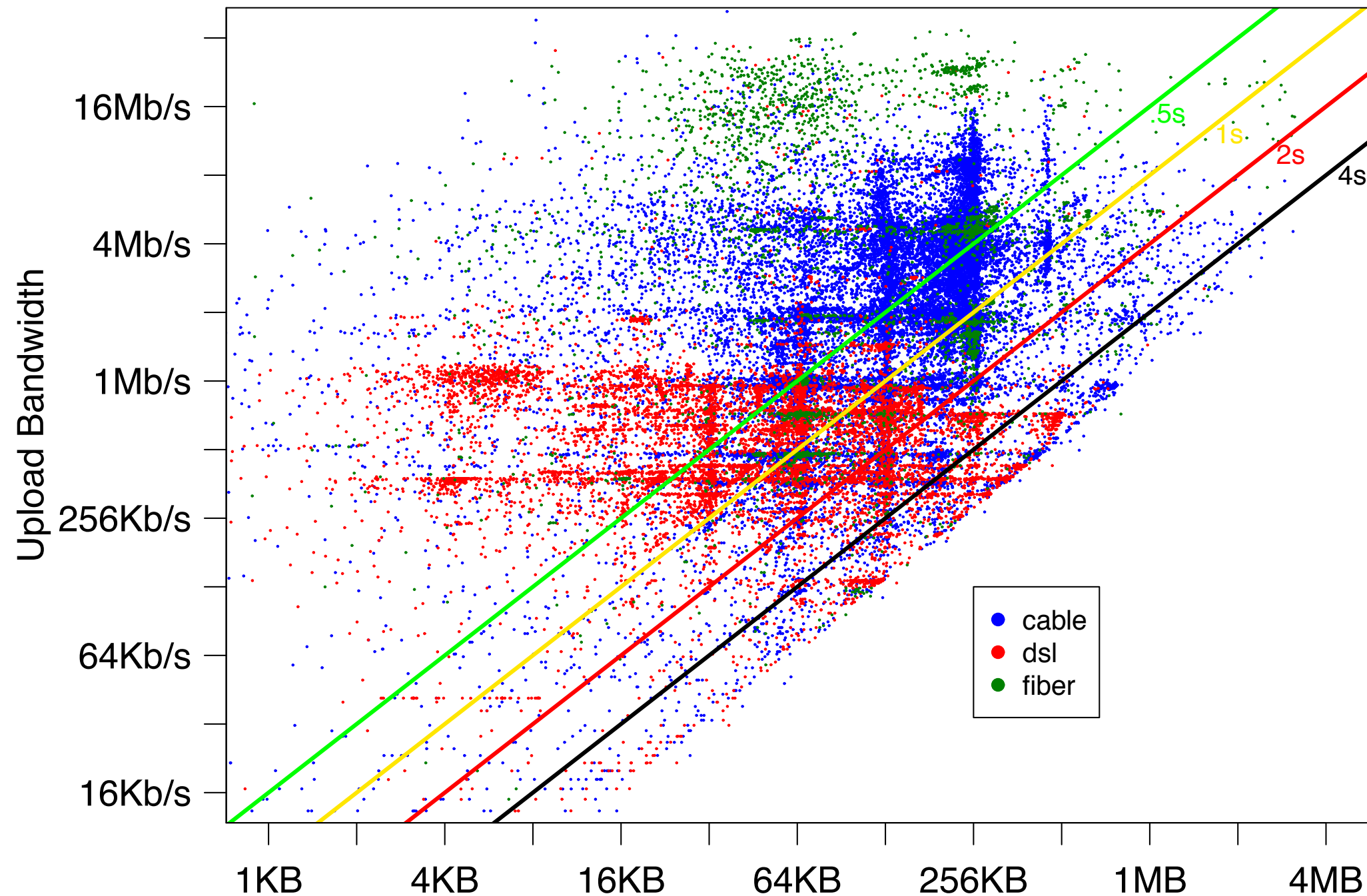
# Network Layer



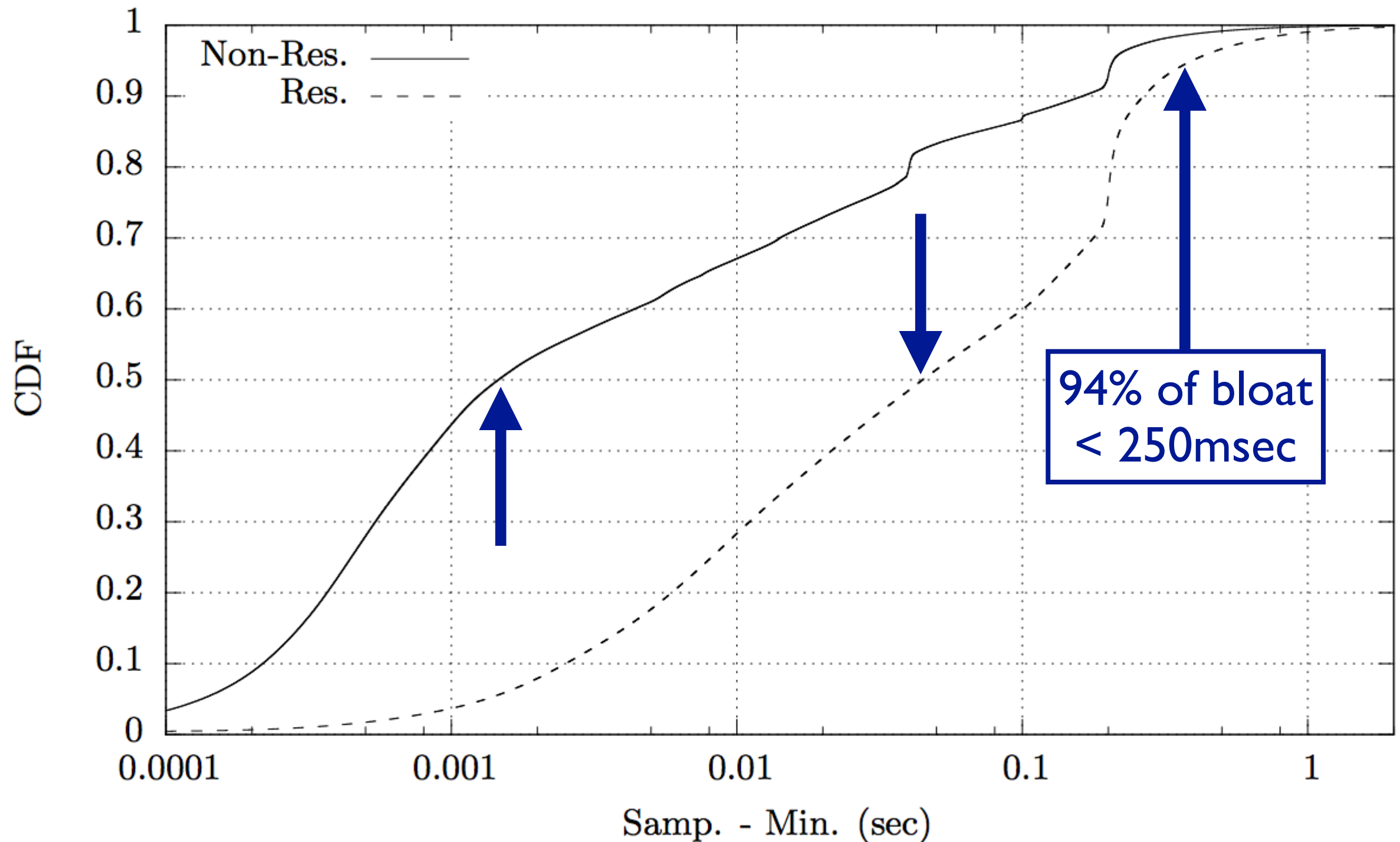
# Bufferbloat

- Claim: over buffering in networking gear is a *big problem*
- “*dark storm clouds surround[ing] us*” (Gettys)
- Case: anecdotal or mis-understood

# Bufferbloat Can Happen ...

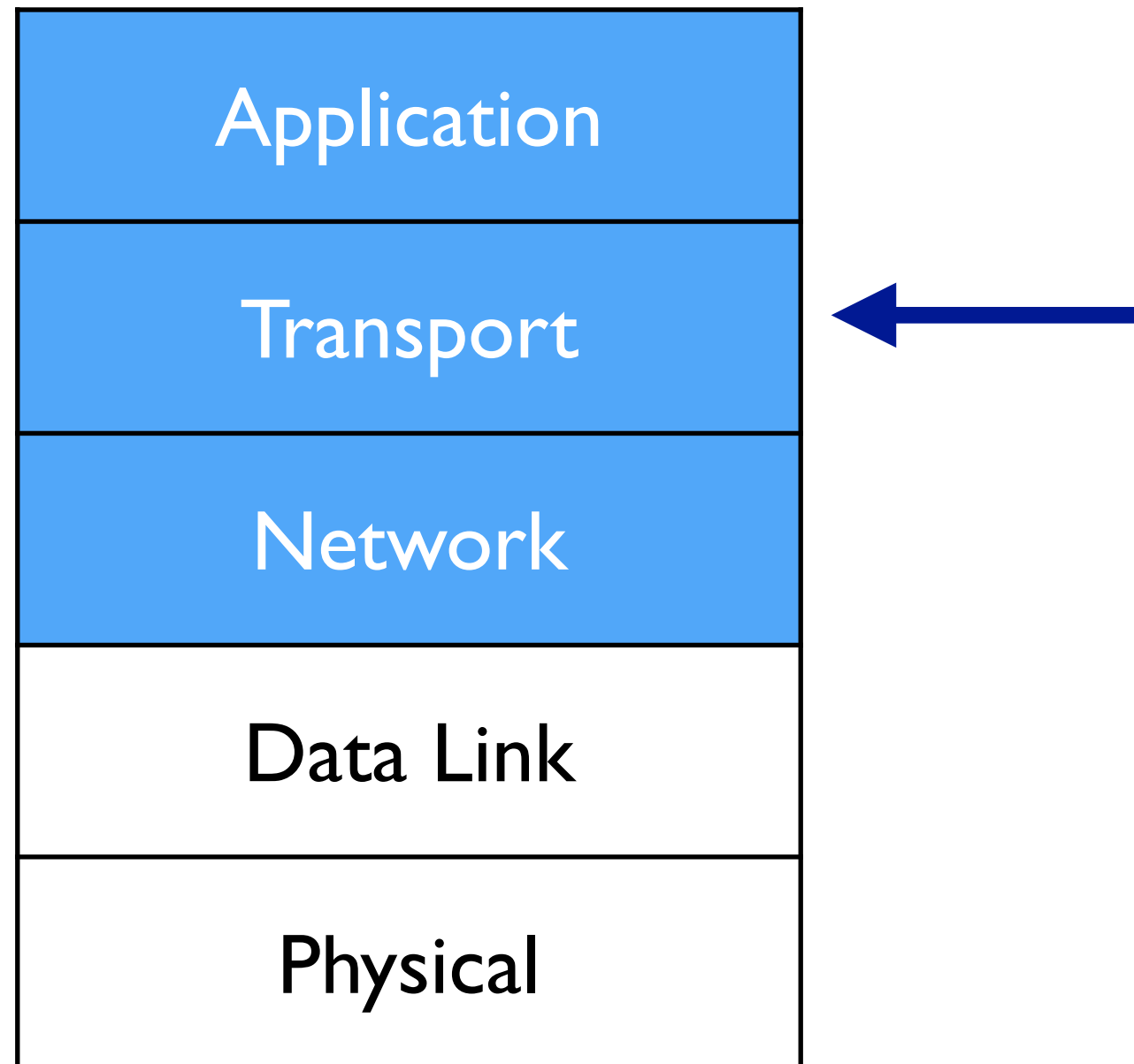


# ... But Doesn't That Much



# It's A Long Way To The Top ...

## Internet Layers



# Transport Layer

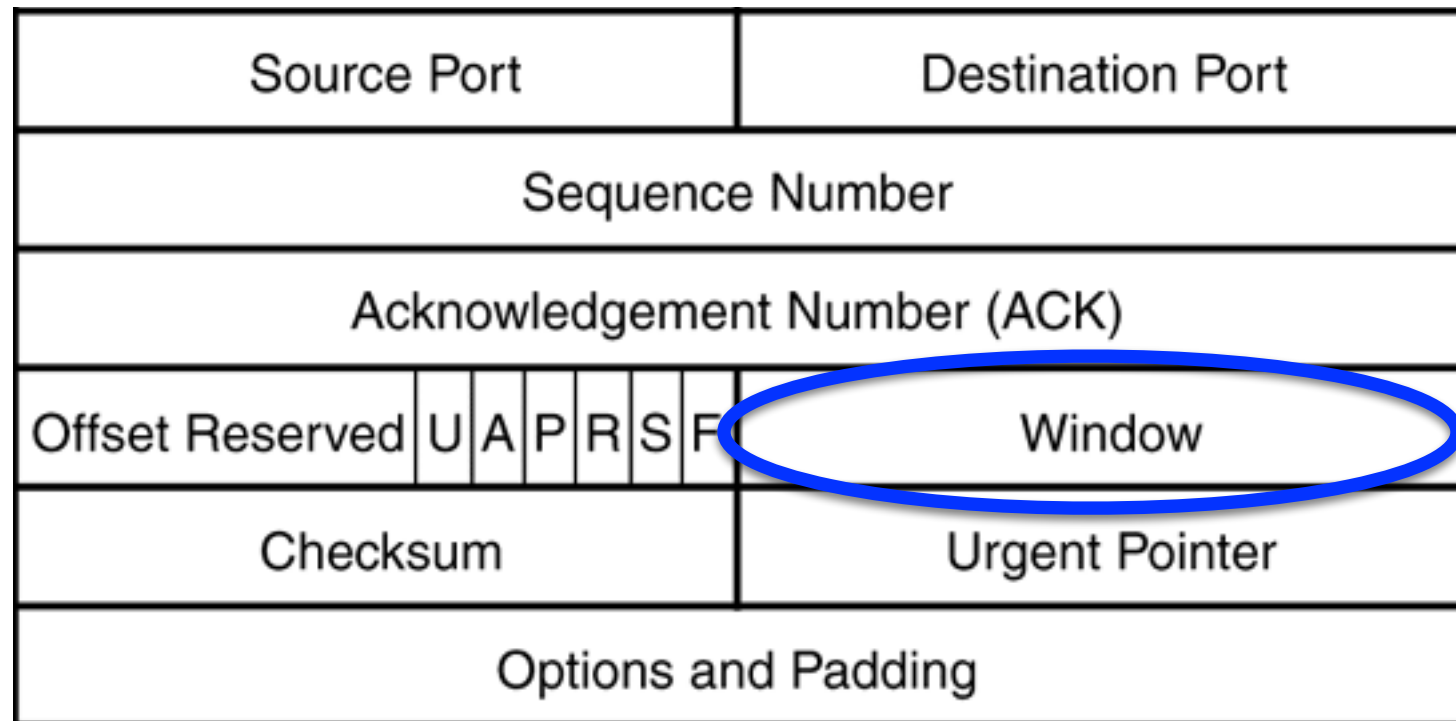
- TCP imposes two kinds of control on traffic
  - flow control
  - congestion control
- Both provide fundamental limits on performance via sliding windows

# Sliding Window Protocols

$$Thput = \frac{Window}{RTT}$$



# Flow Control



$\times 2^{14} = 1 \text{ GB max}$

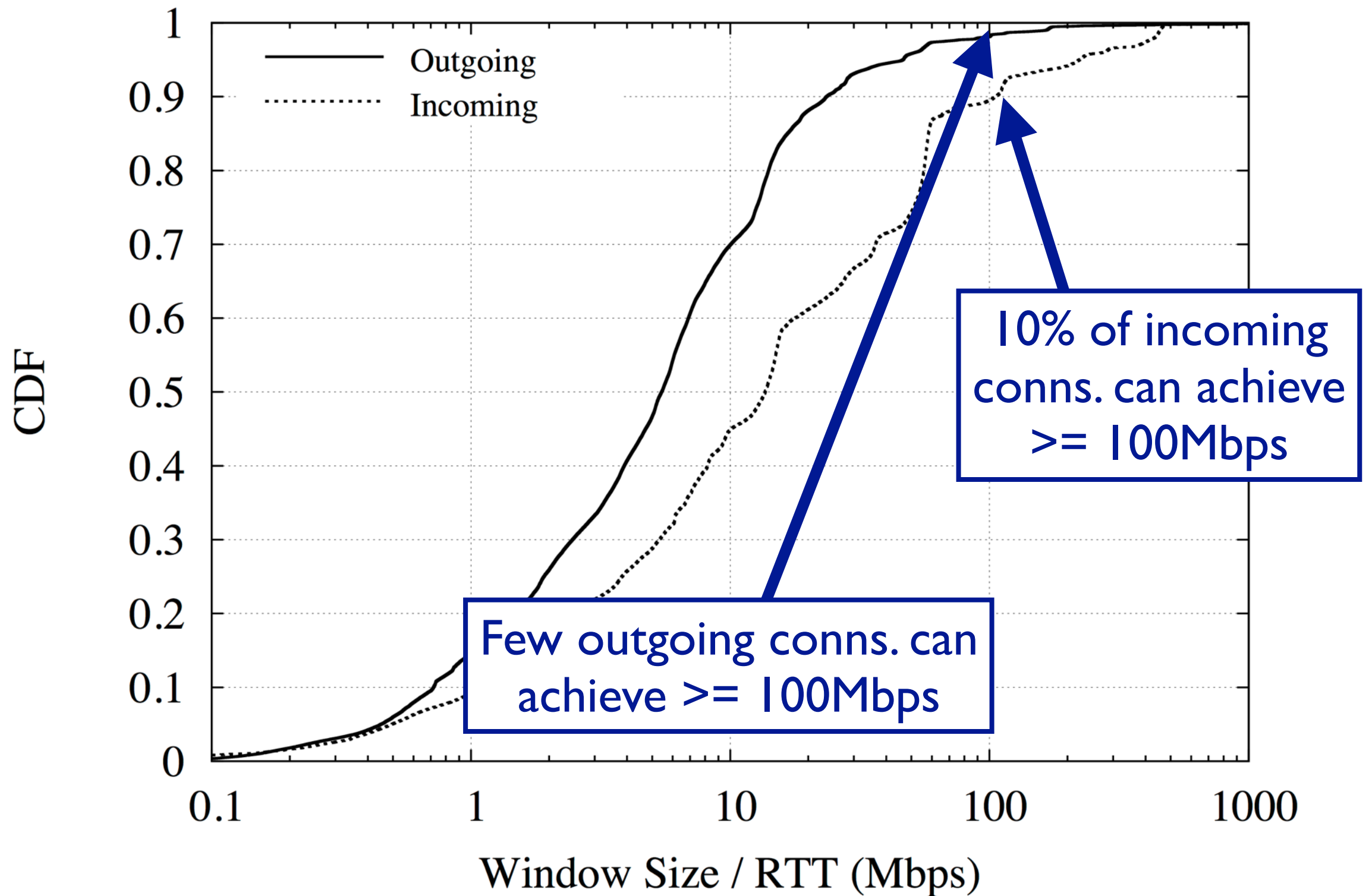


32 bits

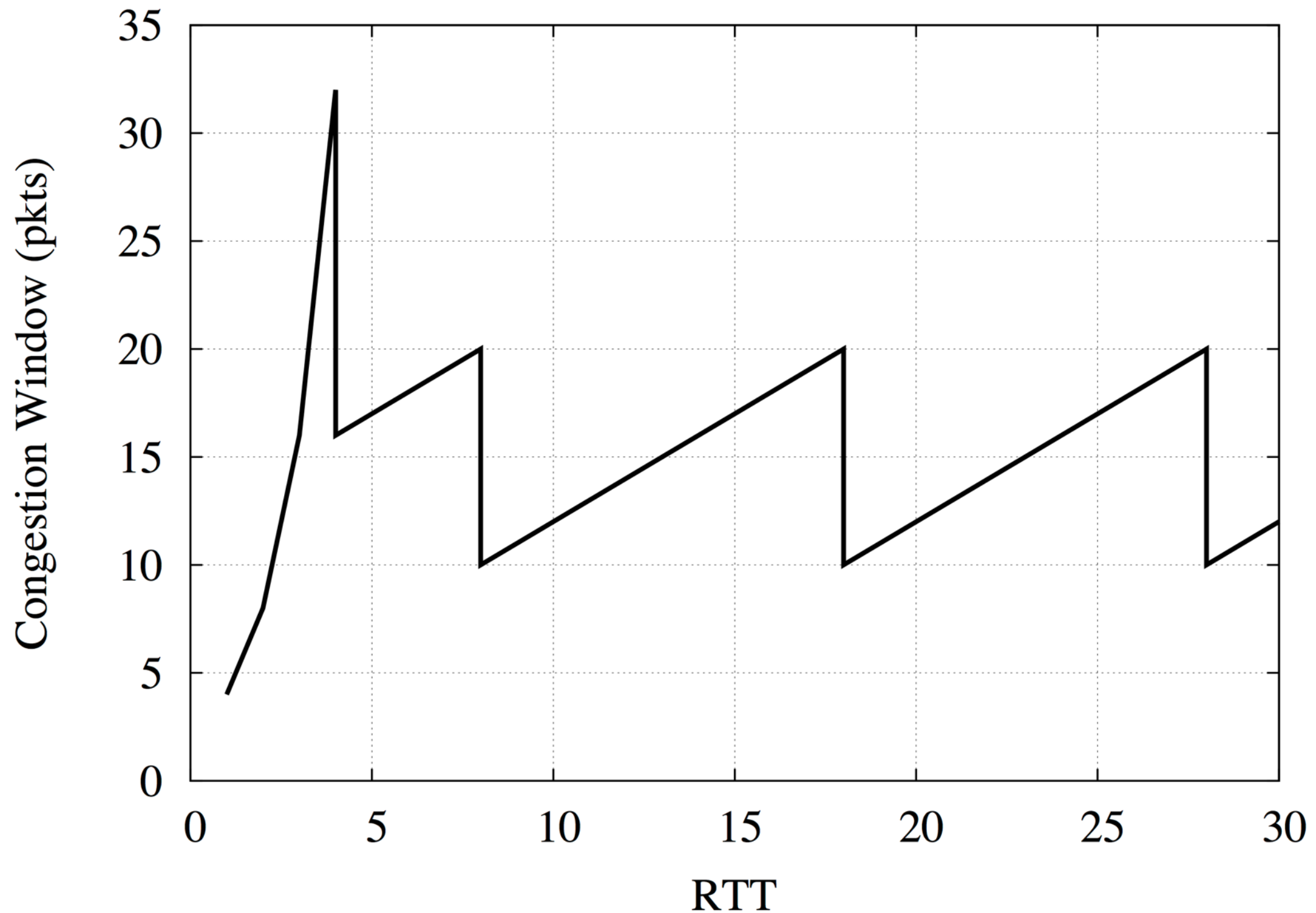
# Flow Control

Window	RTT	Thput
1 MB	1 msec	7.8 Gbps
1 MB	10 msec	800 Mbps
1 MB	100 msec	80 Mbps
64 KB	1 msec	500 Mbps
1 GB	1 msec	8 Tbps
640 KB	500 msec	10 Mbps

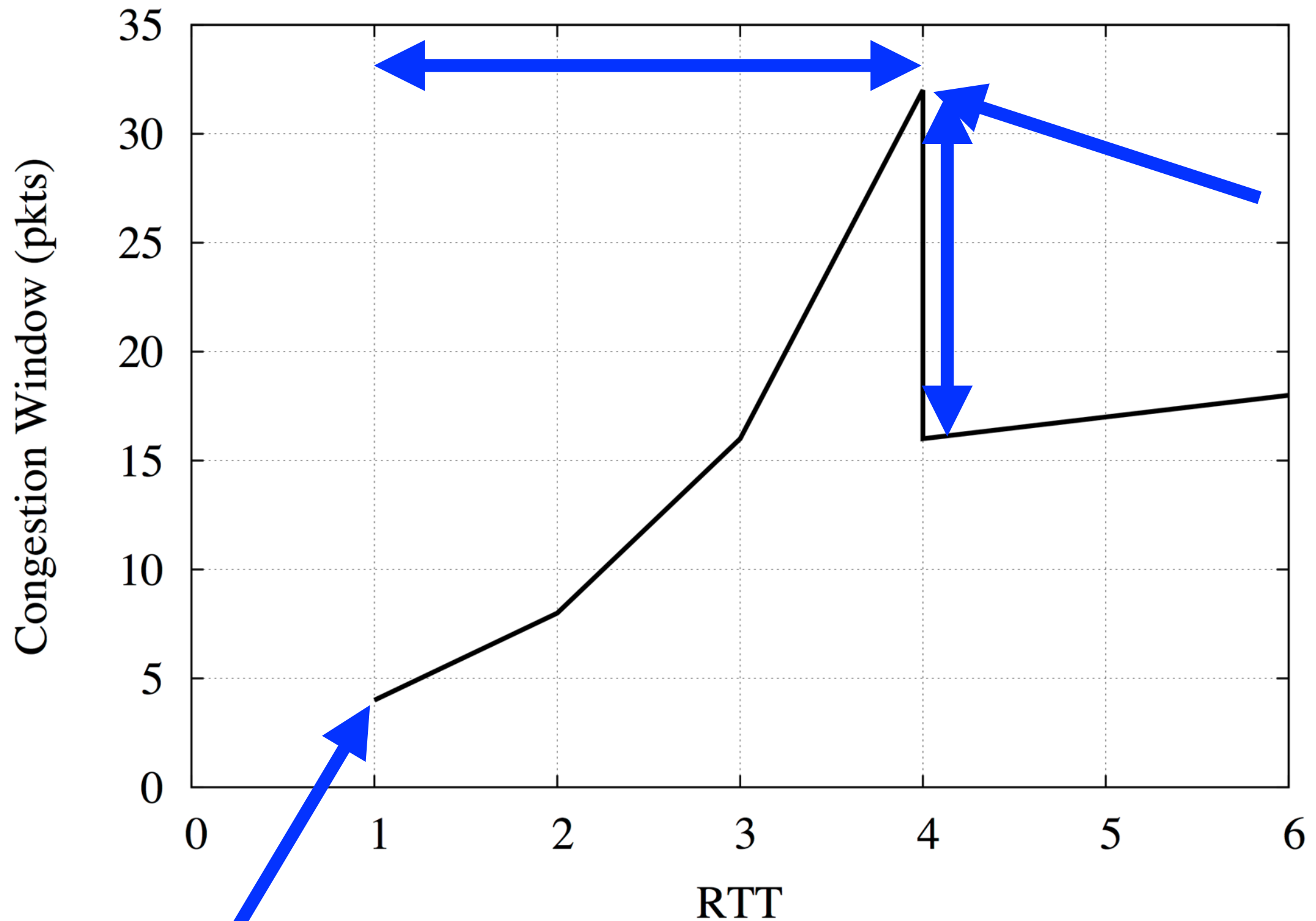
# CCZ Flow Control



# Congestion Control



# Congestion Control

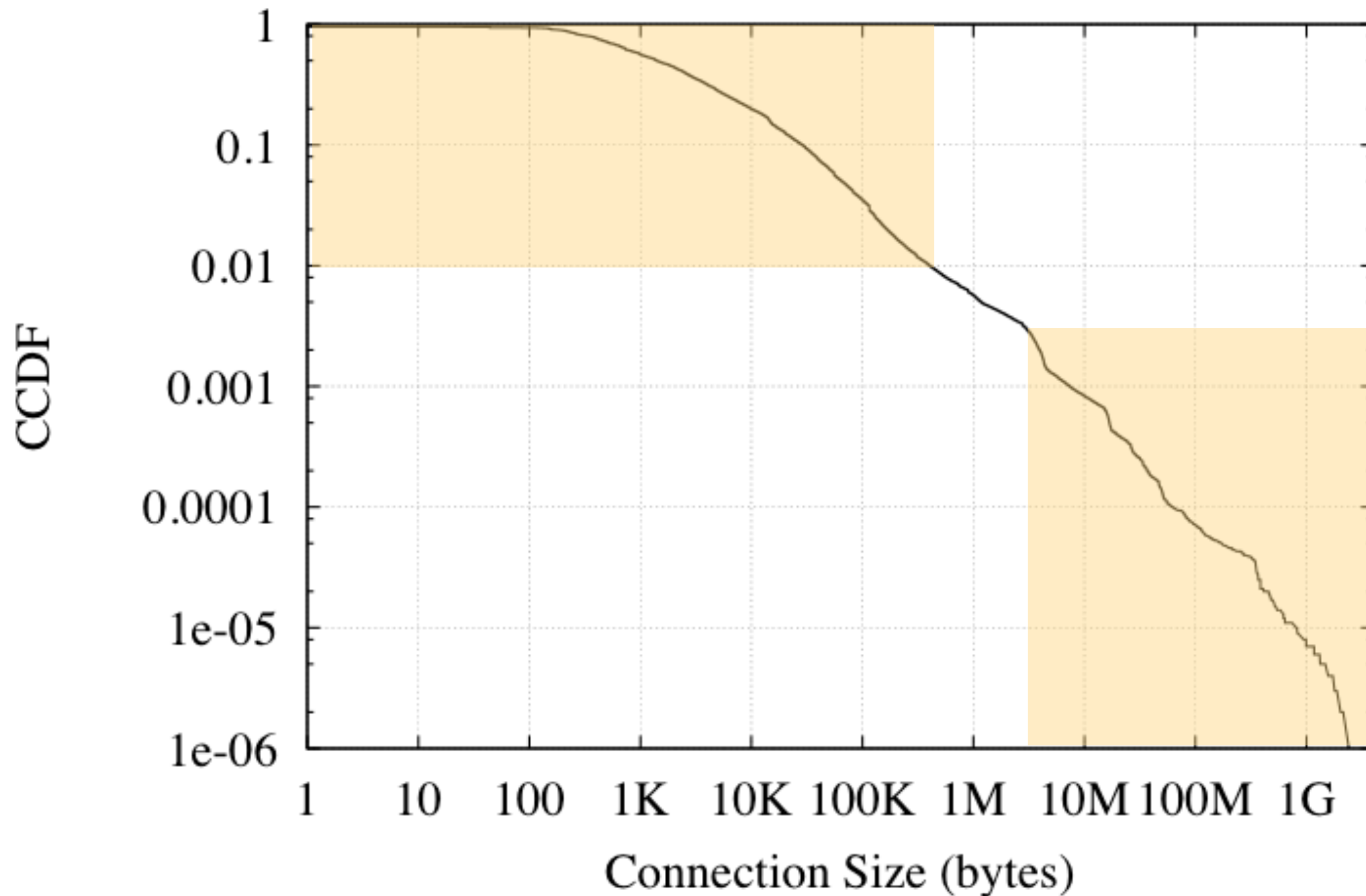


# Congestion Control

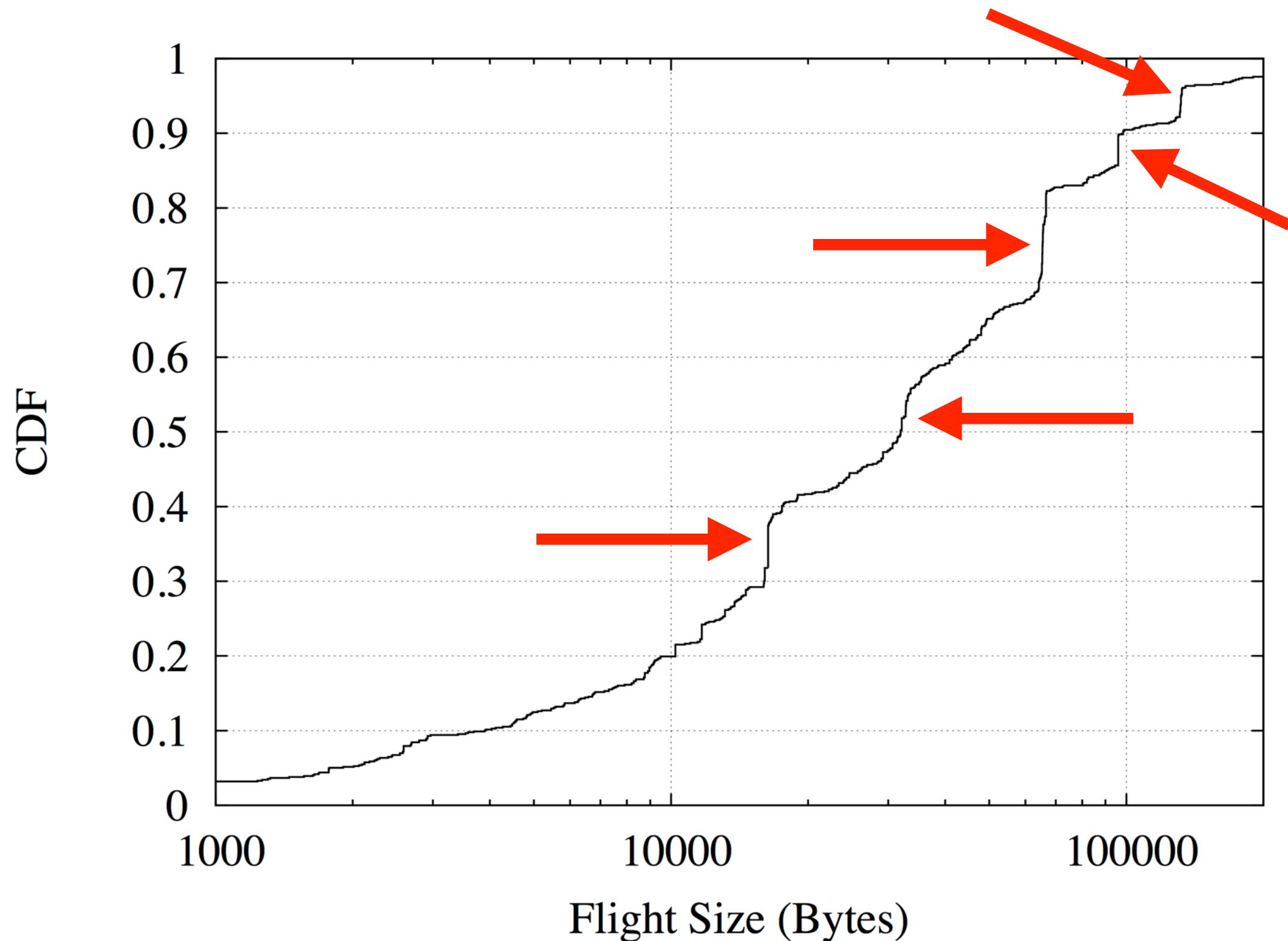
- Takes *data* to build *speed*
- To reach a congestion window of  $X$  bytes requires sending  $\approx 2X$  data bytes

# Connection Sizes

LBNL; Jun 30 2011; 3.4M conns.; 273GB



# CCZ Connections Without Loss

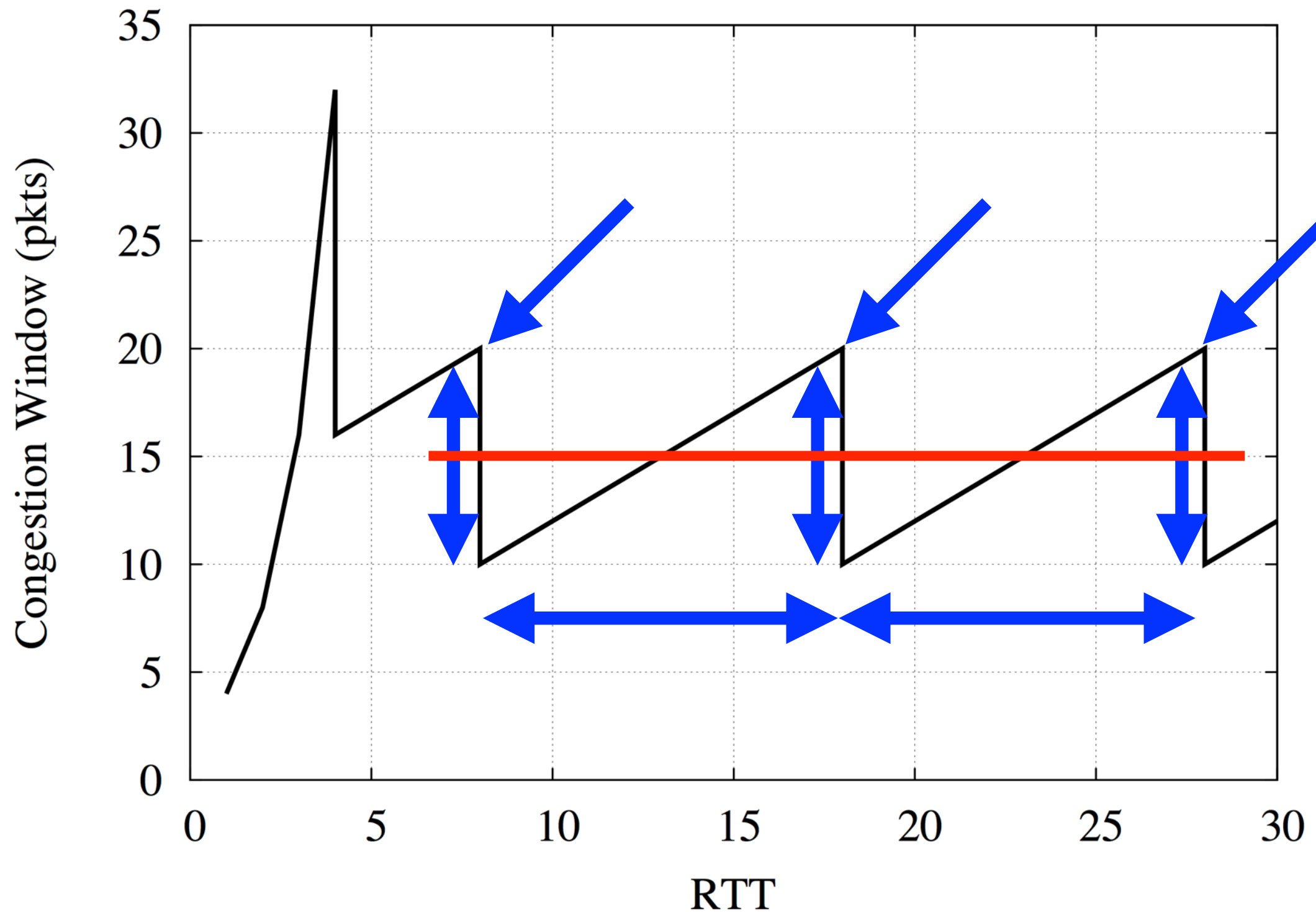




# CCZ Connections Without Loss

- 12% of connections constrained by advertised window
- 45% of connections constrained by a sender-side buffer

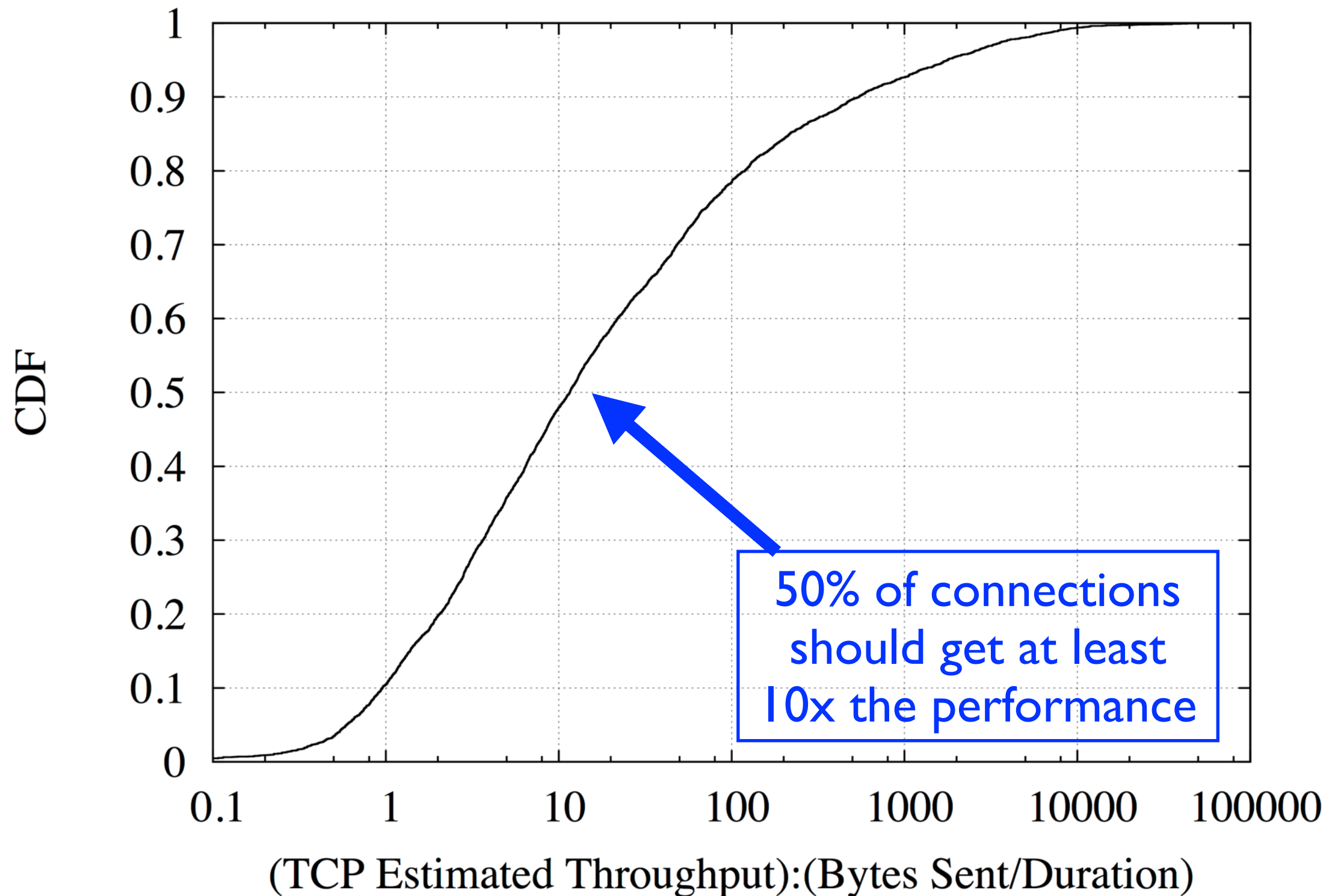
# Congestion Control



# General TCP Performance Model

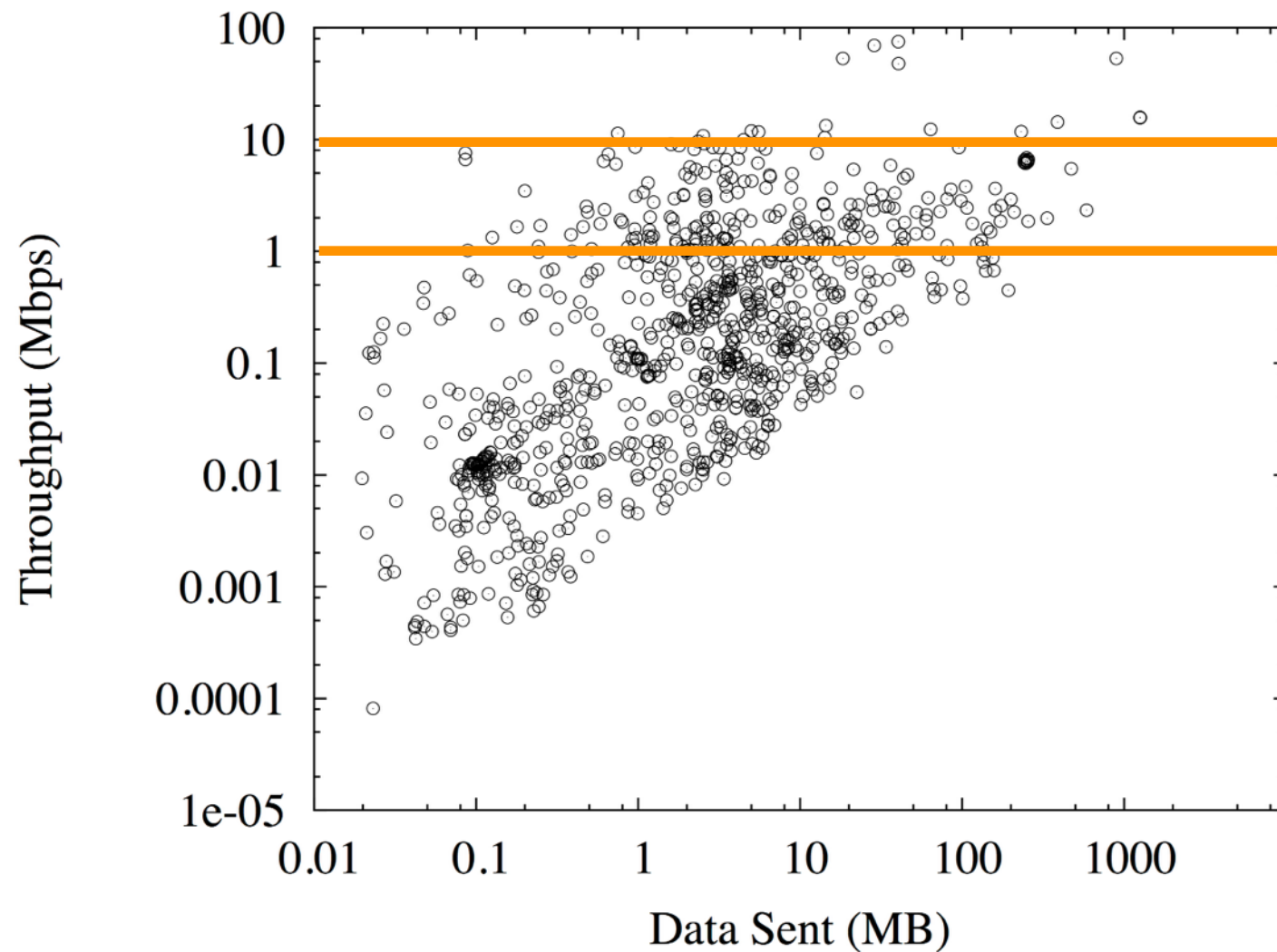
$$BW \propto \frac{MSS}{RTT \times \sqrt{p}}$$

# CCZ Connections With Loss

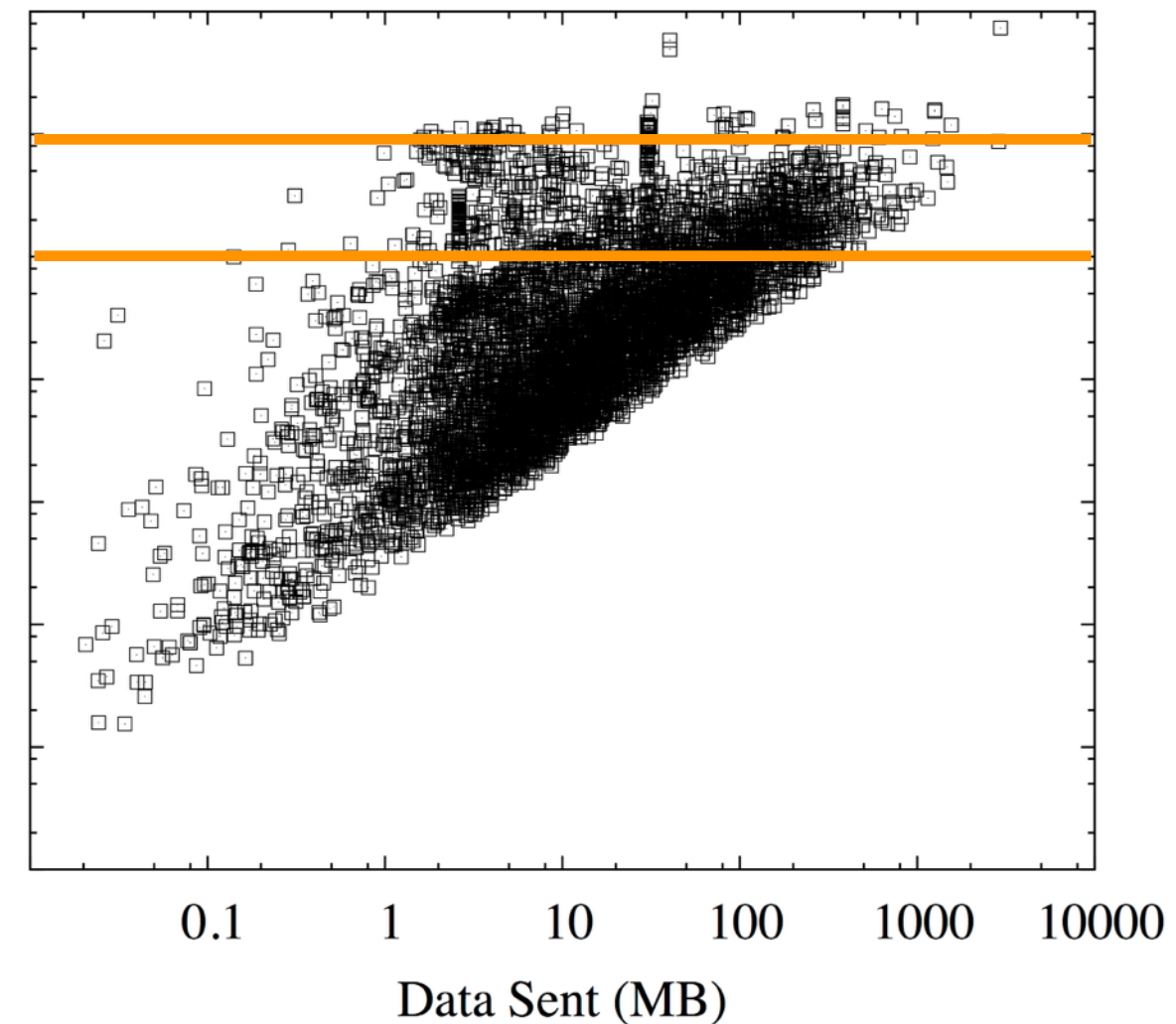


# CCZ Performance

## Connections Without Loss

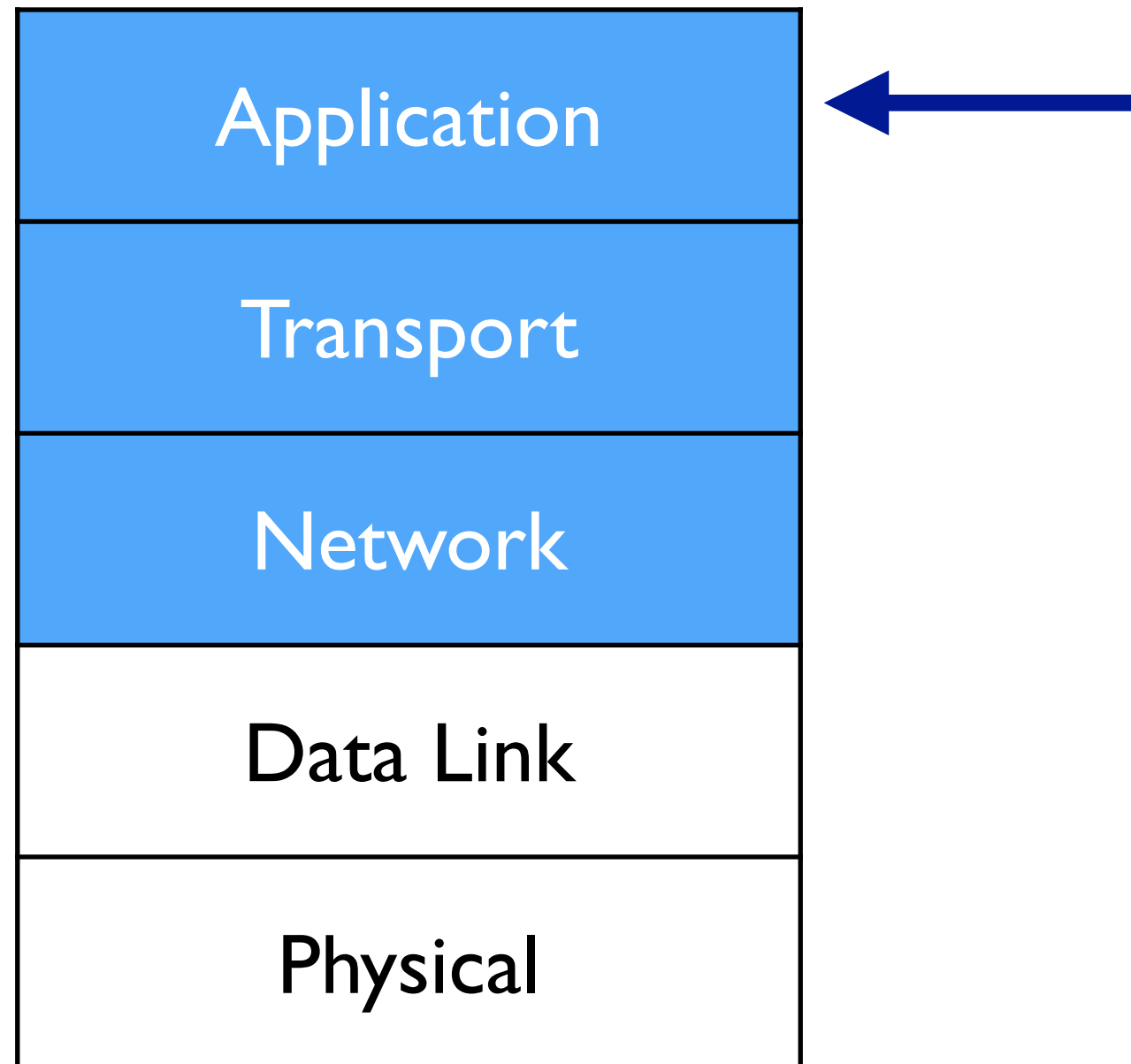


## Connections With Loss

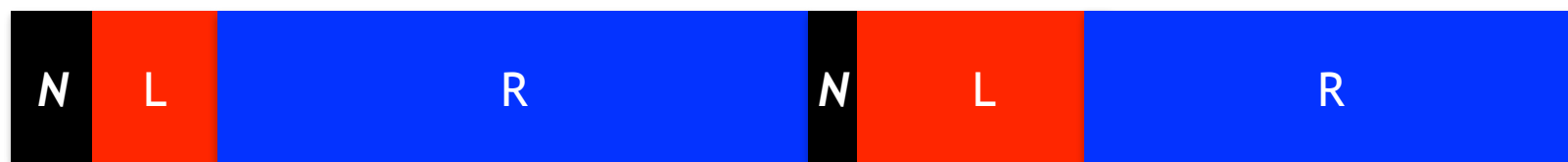


# It's A Long Way To The Top ...

## Internet Layers



# Application Patterns



Time

# Application Use

- We find internal silence in ...
  - ... 28% of connections at ICSI border
  - ... 37% of connections at CCZ border



# Thinking About Capacity

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Application
Transport
Network
Data Link
Physical



# Acknowledgments

- TCP performance work from Matt Sargent (Case EECS)
- Work funded in part by NSF
- Data sources: Case CCZ, LBNL, ICSI



# Questions? Comments?



## References:

*<http://www.icir.org/mallman/perf.html>*



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*<http://www.icir.org/mallman/>*