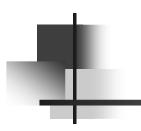
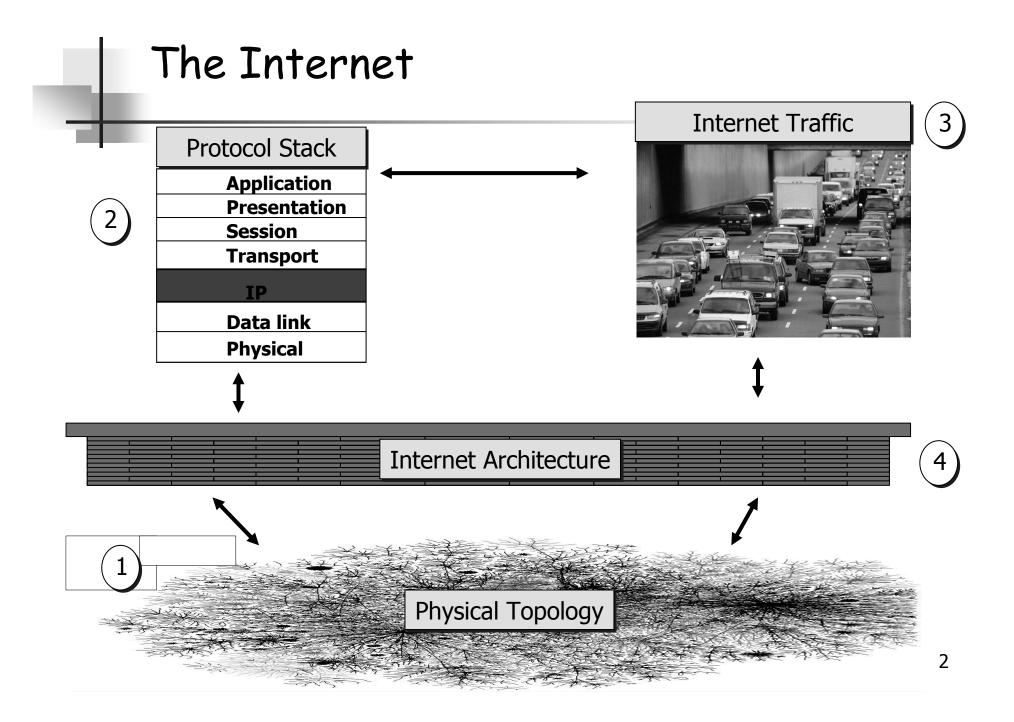
Measuring the Evolution of Transport Protocols in the Internet



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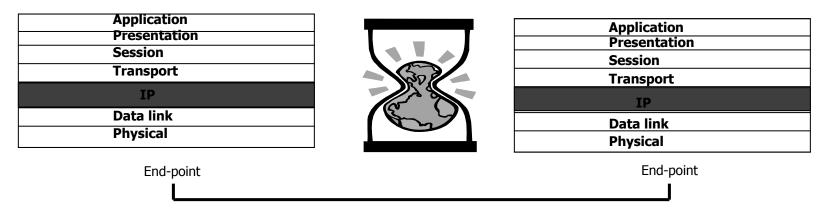


TCP Evolution

- Congestion Control behavior
 - No proper congestion control => congestion collapse
- Deployment, correctness of transport mechanisms
 - Assess correctness and behavior of newer additions
- Dynamics: Theory vs. Practice
 - Differences between <u>protocol specs</u> (theory) and their <u>implementation and its environment</u> (practice)
- Network Modeling
 - Aim at improving accuracy of network models

Network Evolution

Hourglass Model

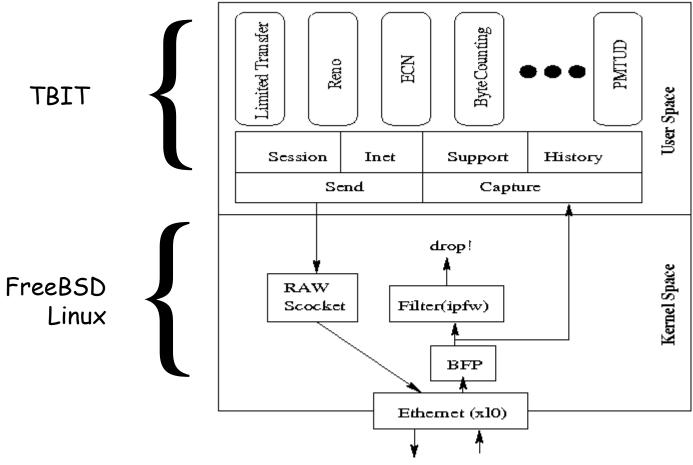


- End-to-end principle
 - "Some functions can only be implemented completely and correctly end-to-end, with the help of the end points"
- Study effect of middleboxes on these principles
 - firewalls, load balancers, NATs, ...

Experimental Platform

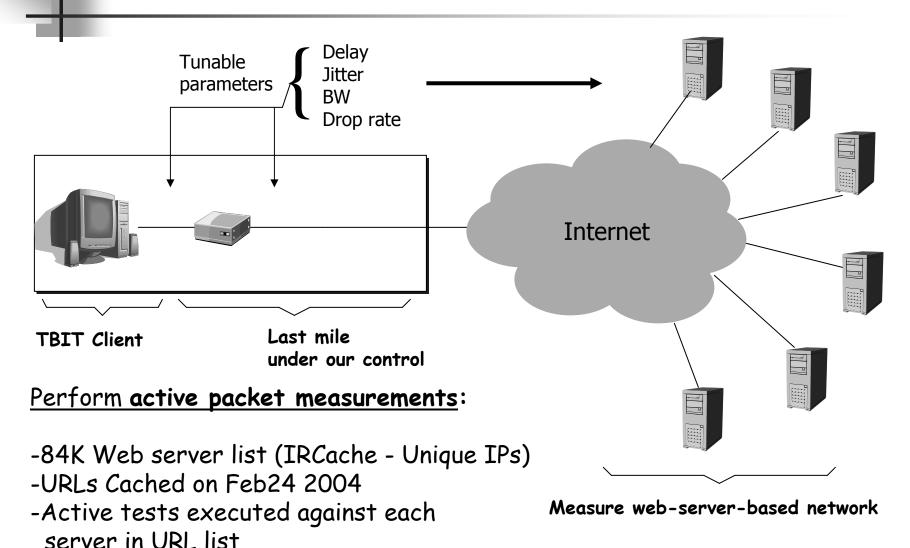
- Measuring TCP implementations
 - <u>Passive measurements</u> of web clients
 - Active measurements
 - Web server mechanisms
 - Interactions with environment
- Active measurements requirements
 - Measure in-the-field Web servers
 - Employ only conformant TCP traffic
 - Unilateral control at measurement side
- Employ "undercover" web clients...

Undercover Web Clients: TBIT

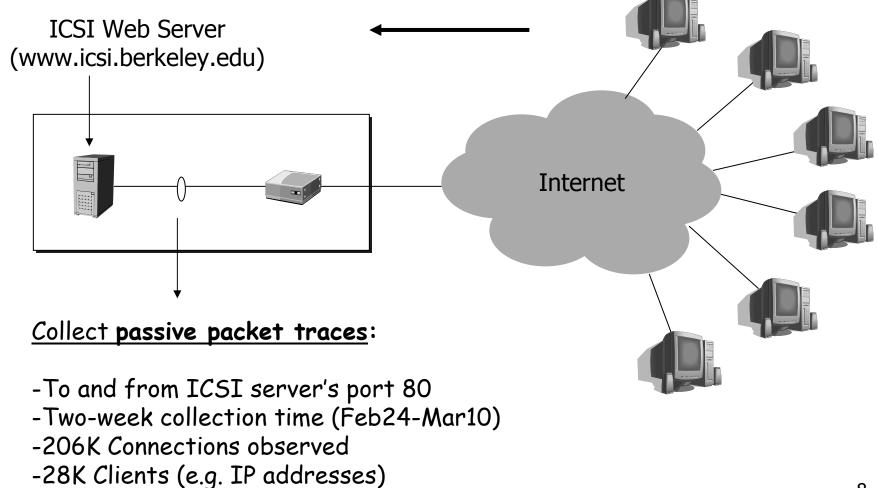


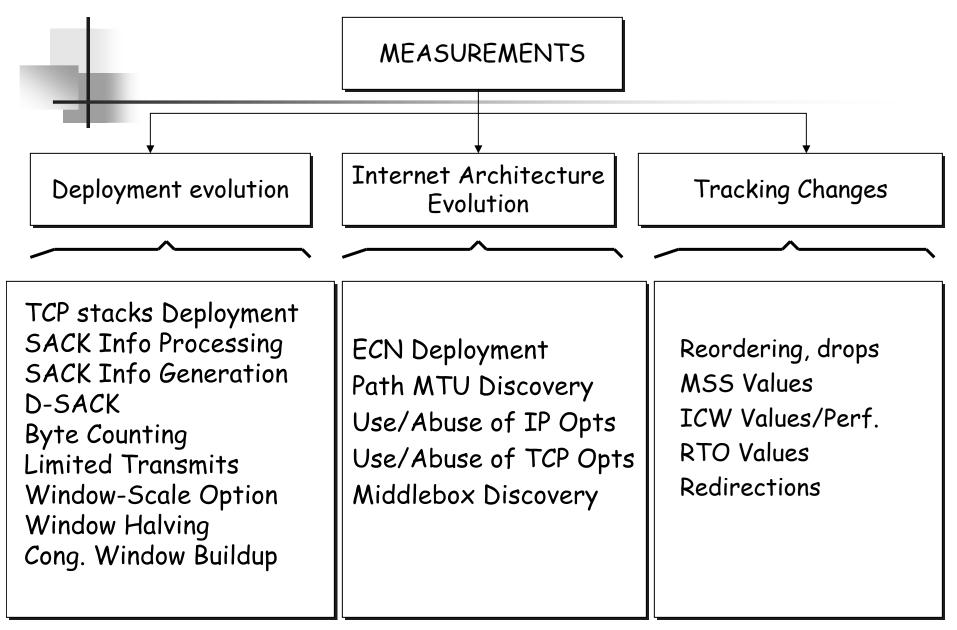
To Web Servers (Internet)

Experimental Platform: Server Side



Experimental Platform: Client Side





Talk Outline

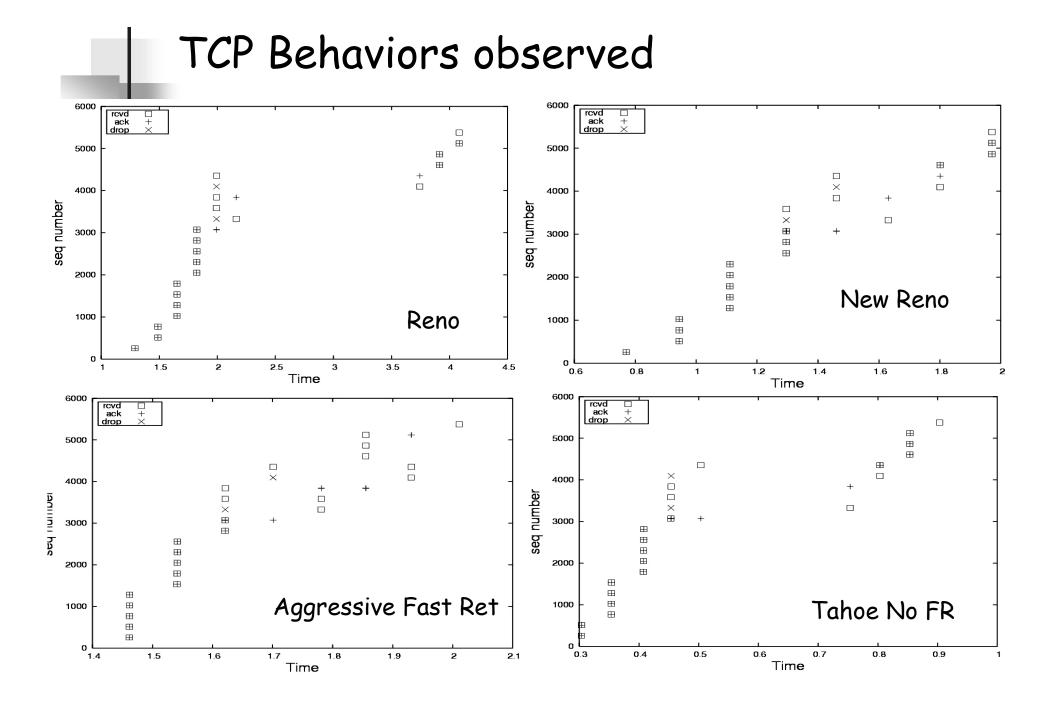
- \checkmark Motivation
- ✓ Measurement Platform
- Active Measurements
 - Deployment of Transport Mechanisms
 - Middleboxes and Transport Protocols
- Summary of Results
 - Including client-side
- Conclusions
- Future Work

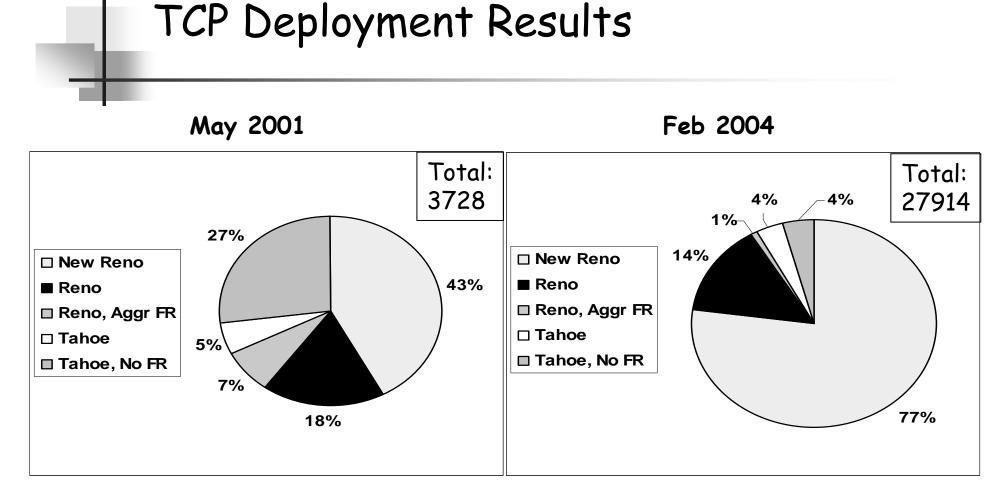


Deployment of Transport Mechanisms

Test: Assess Deployment of TCP stacks

- Establish connection with Web server
 - Use small MSS
 - Restrict Congestion window to 5 segments
- Request web page
- Receive and ACK incoming packets, but...
 - Drop packet 13th
 - Receive and ACK packets 14th and 15th
 - Drop Packet 16th
 - Continue download until receiving packet 25th

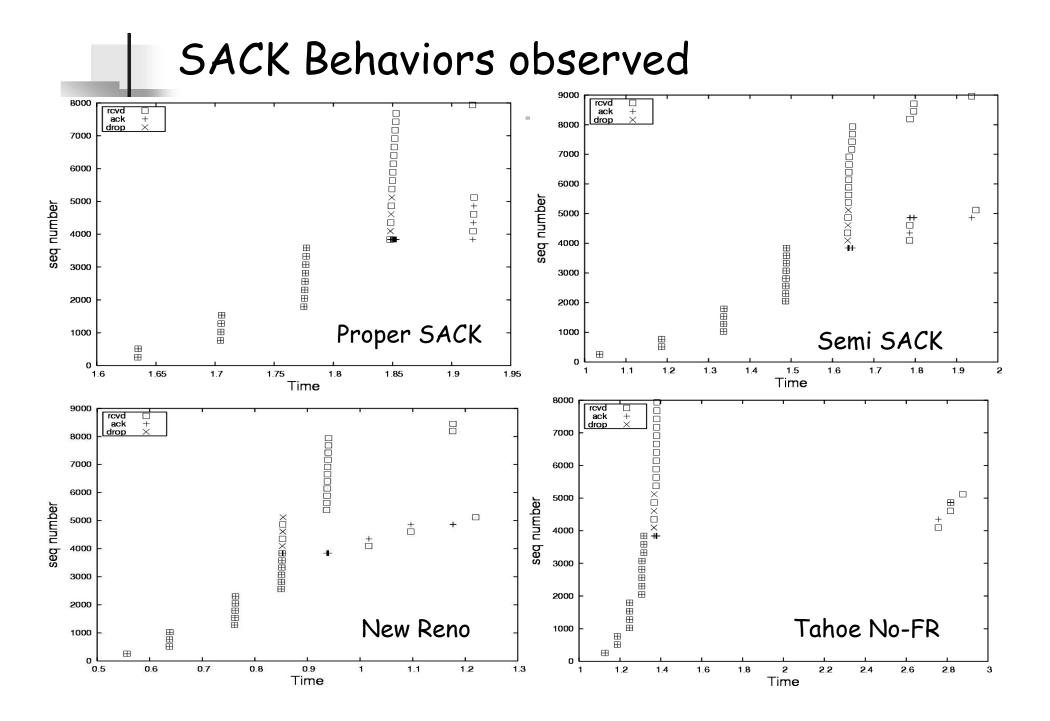


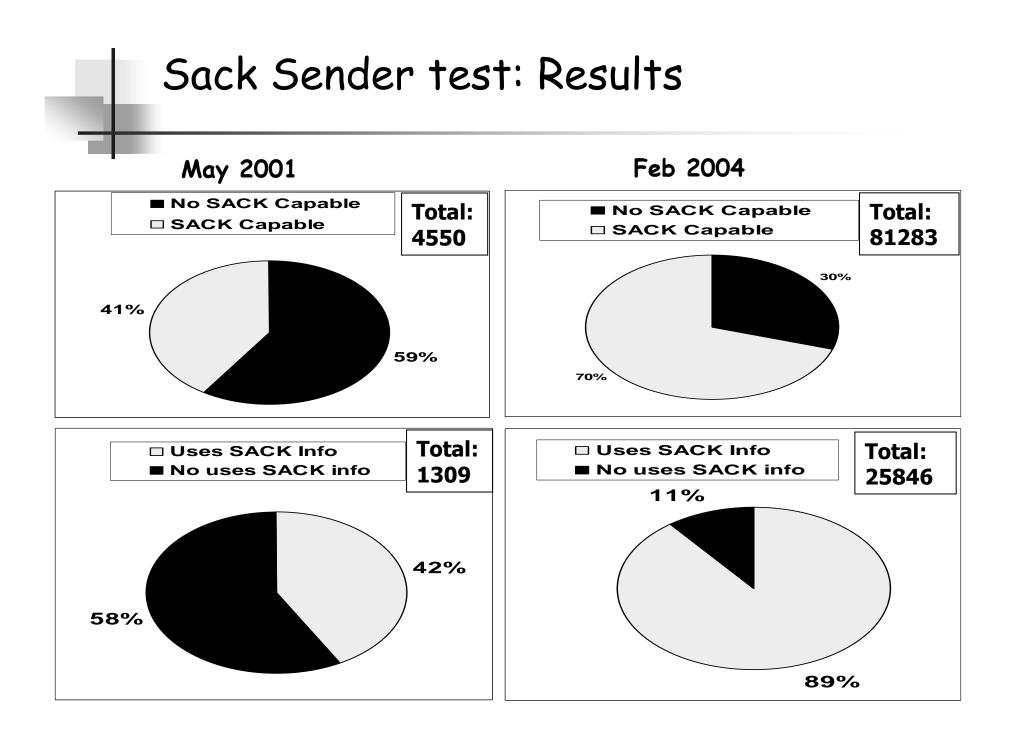


- Deployment of New Reno increased significantly
- Buggy Tahoe without Fast Retransmit decreased
- Network simulations should use New Reno TCP

Test: Assesing SACK Behaviors

- Negotiate SACK-enabled connection
- If server not capable => NO SACK
- Request web page download
- Receive and ACK incoming packets but...
 - Drop packets 15th, 17th, and 19th
 - Continue receiving and ACKing packets normally (sending appropriate <u>SACK blocks</u> for "drops")
 - Observe retransmission behavior
 - Terminate test, close connection





Generation of SACK Information

Do servers generate accurate SACK information?

Т

Test:

Request: "GET / HTTP 1.1... ↓ G X T X / X H X T X \b X . .

"Drop" X-marked packets and update sequence numbers appropriately

	SYN (1)►
	← SYN/ACK
	► ACK►
	→ REQ(`G' (2))→
	< ACK(3)
	→ REQ('T' (4))→
	← ACK(3, SACK(3))
	→ REQ(`/´(6))→
	<pre> ACK(3, SACK(3), SACK(5))</pre>
B	SIT Serve

Sack Receiver Test: Results

Type of Server	Servers (%)
Total servers	84394
I. Not SACK-Capable	24361 (28.8%)
II. SACK blocks OK	54650 (64.7%)
III. Shifted SACK blocks	346 (0.5%)
I.V. Errors	5037 (6%)

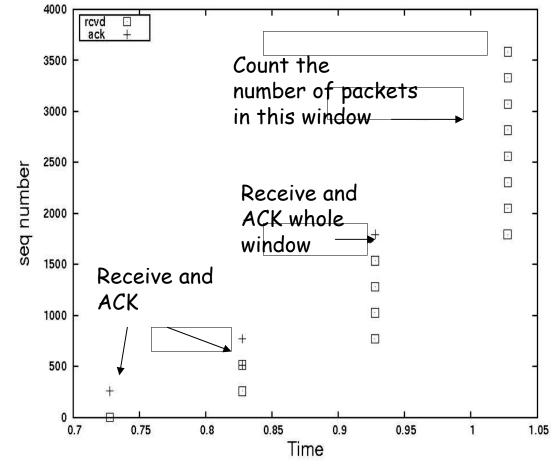
- Shifting blocks could have been caused by:
 - NATs, Fingerprint scrubbers,...
- Such middlebox interactions affect any TCP-based communication

Test: Appropriate Byte Counting (ABC)

- TCP Congestion Control
 - Slow start: increase CWND by one segment for each received ACK
 - Congestion avoidance: increase CWND by 1/MSS for each received ACK
- Drawbacks
 - Delayed ACKs reduce CWND opening rate
 - Mis-behaving receivers may induce servers to open CWND too fast
- ABC Proposal
 - Increase CWND based on bytes ACKed by incoming ACKs, instead of based on number of ACKs received

ABC Test: Example for ICW = 1

- Receive and ACK packets 1, 2 and 3
- Wait for window of 4 packets to arrive
- ACK whole window
- Count number of packets received in next window



ABC Test: Results

Slow Start Behavior	Number (%)
Total number of servers	44579
I. Classified Servers	23170 (52%)
I.A. Packet Counting	15331 (51.9%)
I.B. ABC	65 (0.1%)
II. Unknown behavior	288 (0.6%)
III. Errors	21121 (47%)

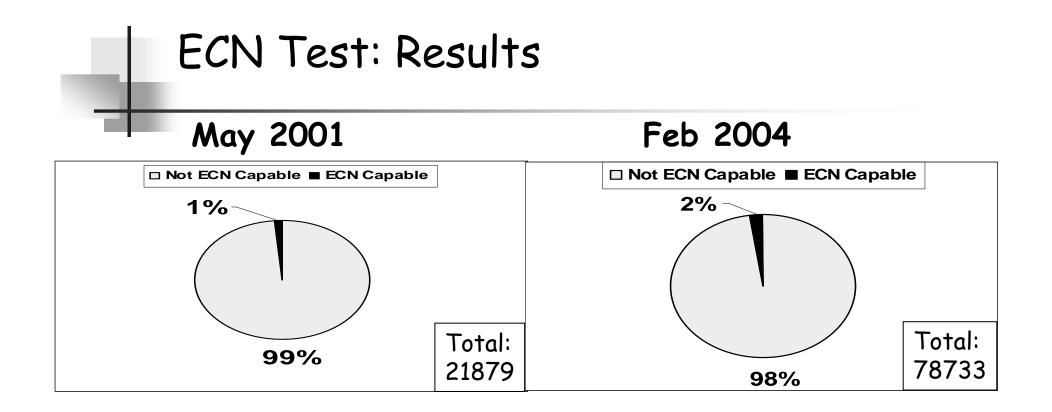
 Notice a <u>5-year old</u> proposed mechanism addressing (1) <u>performance</u> concerns and (2) <u>security</u> issues and yet, not being deployed!

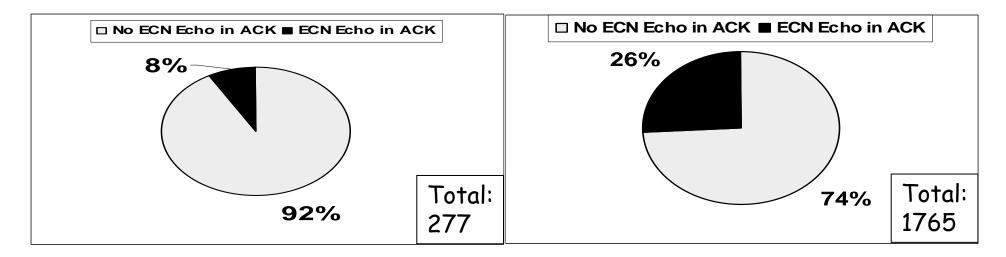


Middleboxes and Transport Protocols

ECN Capabilities

- ECN: Explicit Congestion Notification
 - Allows routers to notify congestion to end nodes
- TCP: 2-way handshake
 - Active end: send ECN-Setup SYN (ECN_ECHO, CWR)
 - Passive end: send ECN-Setup SYN/ACK (ECN_ECHO)
- IP: 2-bit ECN field in IP header => 4 ECN CPs
 - 00: Not ECT
 - 01: ECT(1) Sender is ECN capable
 - 10: ECT(0) Sender is ECN capable
 - 11: Congestion Experienced (CE)

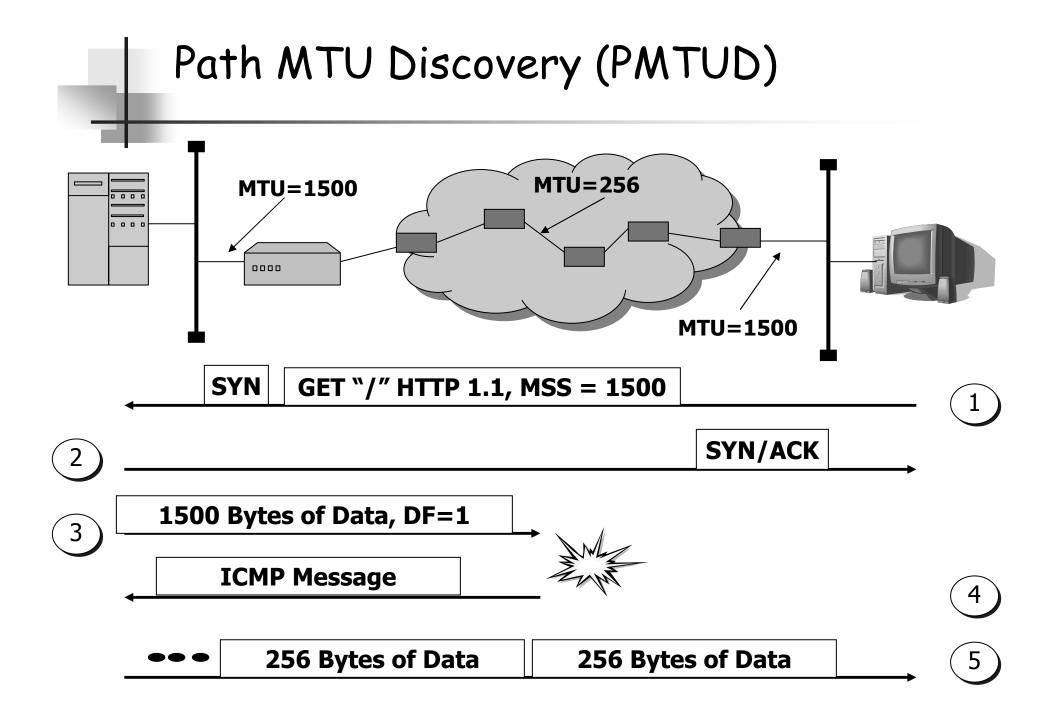


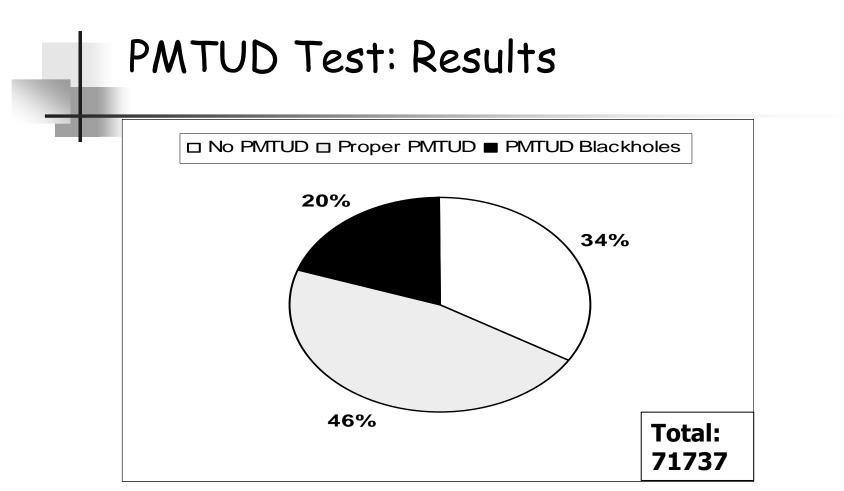




Blocking ECN (3194 Conns) ECN Code Points (1765)

■ With ECN □ Without ECN	ECN CPs data pkts	Number (%)
250/	Received pkts w/ ECT 00 (Non-ECT)	758 (42%)
	Received pkts w/ ECT 01 (ECT(1))	0 (0%)
	Received pkts w/ ECT 10 (ECT(0))	1167 (66%)
	Received pkts w/ ECT 11 (CE)	0 (0%)
75% Total: 3194	Received both pkts with ECTOO & ECT 10	174 (10%)

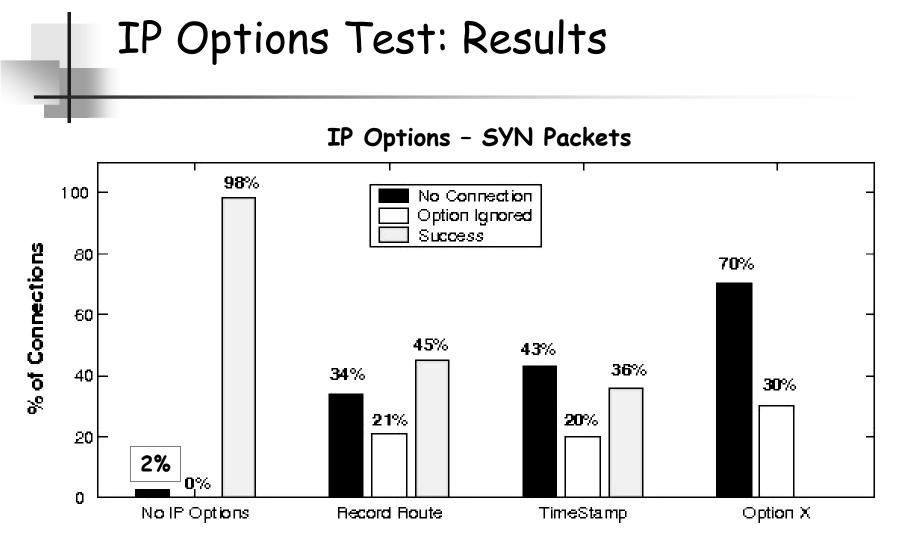




- Observed a non trivial number of black holes
- No hope for new ICMP-based proposed mechanisms
 - Explicit corruption notification
 - Handoff notification

Interference with TCP/IP Options

- TCP/IP options
 - Allow encoding additional information at end of packets
- Several concerns raised about using IP Options
 - Overhead, misalignment problems, DoS attacks
- Solutions to concerns
 - Range from OS patches to dropping "offending" packets
- Issue concerns protocol designers
 - Use of unused TCP/IP options in new proposals
 - Ex: QuickStart (QSR) IP Option
- TCP/IP Options tests
 - Evaluate connections with SYN-packet TCP/IP options
 - Evaluate connections with Mid-Stream TCP/IP options



- Severe interference with known and unknown IP options
- Negative results for of new IP-option-based mechanisms

Summary of Results: TCP Evolution

TCP Mechanism	Conclusion
TCP Cong. Ctrl	~ 2/3 use <u>NewReno</u> => Use it in ns
Loss Recovery	<u>SACK-cap Prevalence</u> : ~ 2/3 servers, 9/10 clients Most claiming SACK, do SACK properly <u>SACK info:</u> (mostly) correct
DSACK	~ $\frac{1}{2}$ of SACK-capable servers, send D-SBs
ABC	Not deployed
LT	Not fully deployed (~1/4 of servers)
MSS	Most clients use ~ 1.4K bytes Most servers accept << 1.4K bytes
RTO	Many servers use RTO < 1s

Summary of Results: TCP Evolution (2)

TCP	Conclusion	
Mechanism		
ICW	Many ICW = 1, Most used ICW = 2-4	
	Some gain from larger ICWs	
	No changes for reordering and losses	
Window Scaling	Most servers support WS (shift count=0)	
Window Halving	Most servers do proper window halving	
	Some servers use CWND without caring for RWND	
Window Buildup	Most servers do no increase cwnd if not used	
Advertised Most clients surveyed advertise 64KB windows		
Window	Many clients advertise 8KB and 16KB	
ECN	Very few servers using ECN (~2.3%)	
	1% Increase since 2001	

Summary of Results: Network Evolution

Behavior	Conclusion
SACK	Small number of cases, web servers and clients receive SACK blocks with incorrect sequence numbers
ECN	Roughly 1% of refused connections
PMTUD	< ¹ / ₂ servers PMTUD-capable Likely routers/middleboxes blocking ICMP messages for 1/6 of the servers
IP Options	Many failures (1/3) when IP RR or TS SYN options used Majority of failures (70%) if unknown IP option used
TCP Options	More resilient and tolerant than unknown options
Reordering	Significant small-scale reordering

Conclusions

- Achieved set goals
 - Tracked deployment of transport mechanisms
 - Evaluated transport-network interactions
- Competition of interests complicates deployment priorities
 - ABC not implemented, LT implemented
 - PMTUD failing, no ECN deployment,...
- Pinpointed specific cases exemplifying how evolving network challenges end-to-end principle
 - Fundamental design principles of Internet have changed
 - Current network needs to evolve towards new reality

Future Work

- Further TCP in-the-field behavior
 - Restart behavior after an idle period; Backoff behavior
 - Behavior in other environments (p2p, wireless,...)
- Study other protocols and mechanisms
 - Existing: UDP, FTP, HTTP, RTP, ...
 - New: AQM, High-Speed TCP, SCTP, DCCP,...
 - TCP in other environments (P2P, web caching,...)
- Further exploration of Middlebox behavior/impact
 - Many open questions (e.g. How about PEPs?)
 - Detecting middleboxes
- Continuous Monitoring Platform
- Active measurements of client behaviors
- Unilaterally-controlled Active Measurements

Contact & Information

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