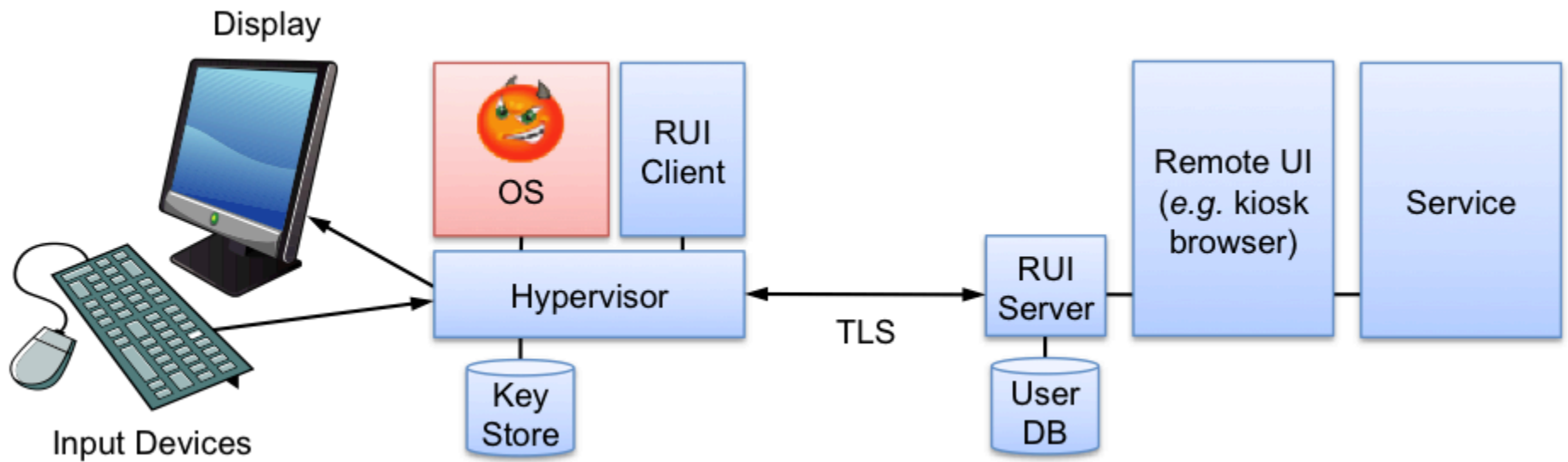
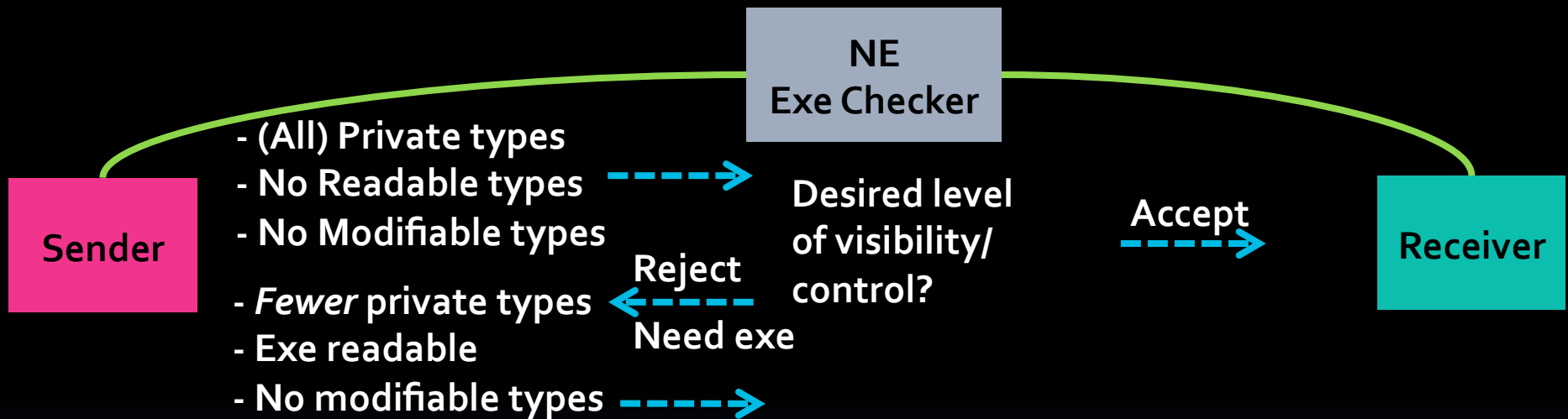


Figure 1: Structure of the Bro system



# Dialog

Typing paves way for *dialog* to negotiate communication properties

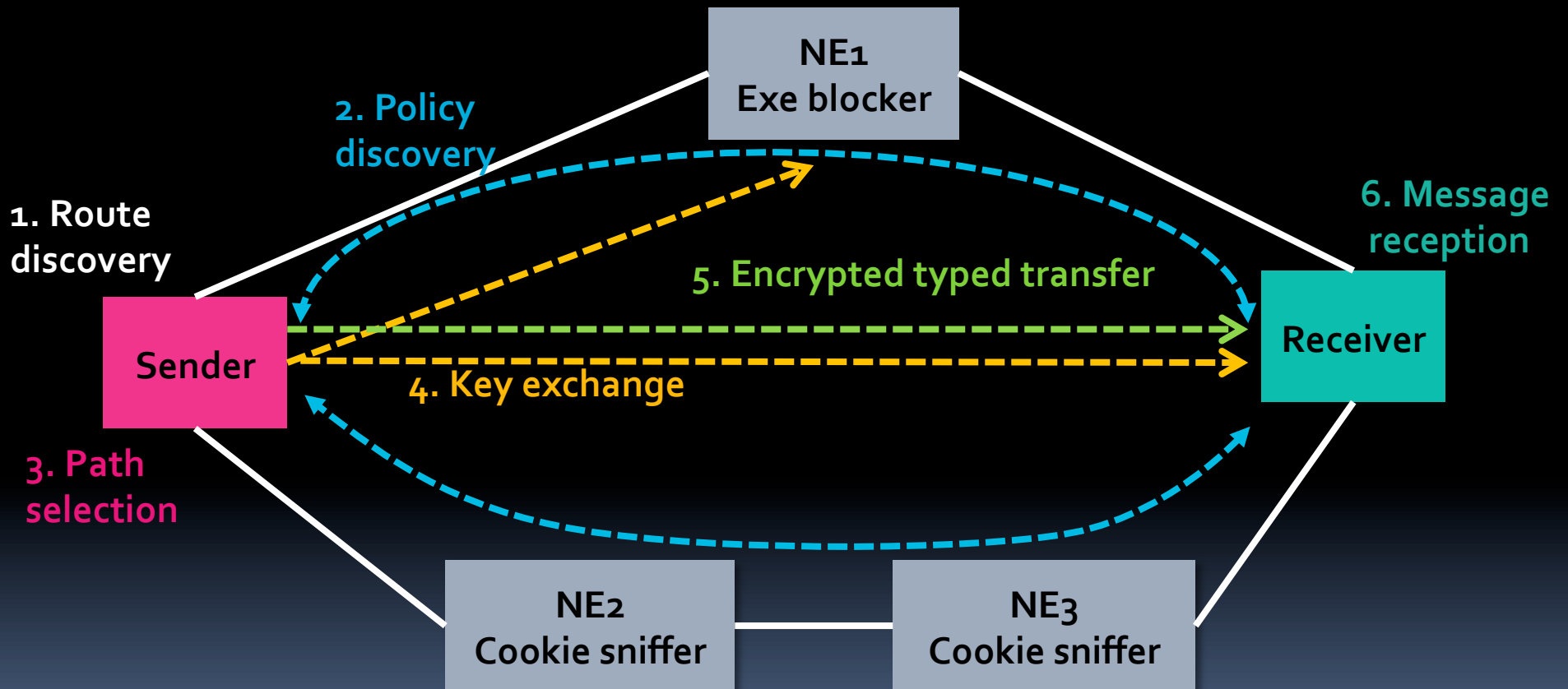


Pre-connection or in-band

Sender may choose an alternate path.  
Fail if no such path → *reason in full view*

Network has upper hand, but visibility limits collateral damage

# Progression of Communication

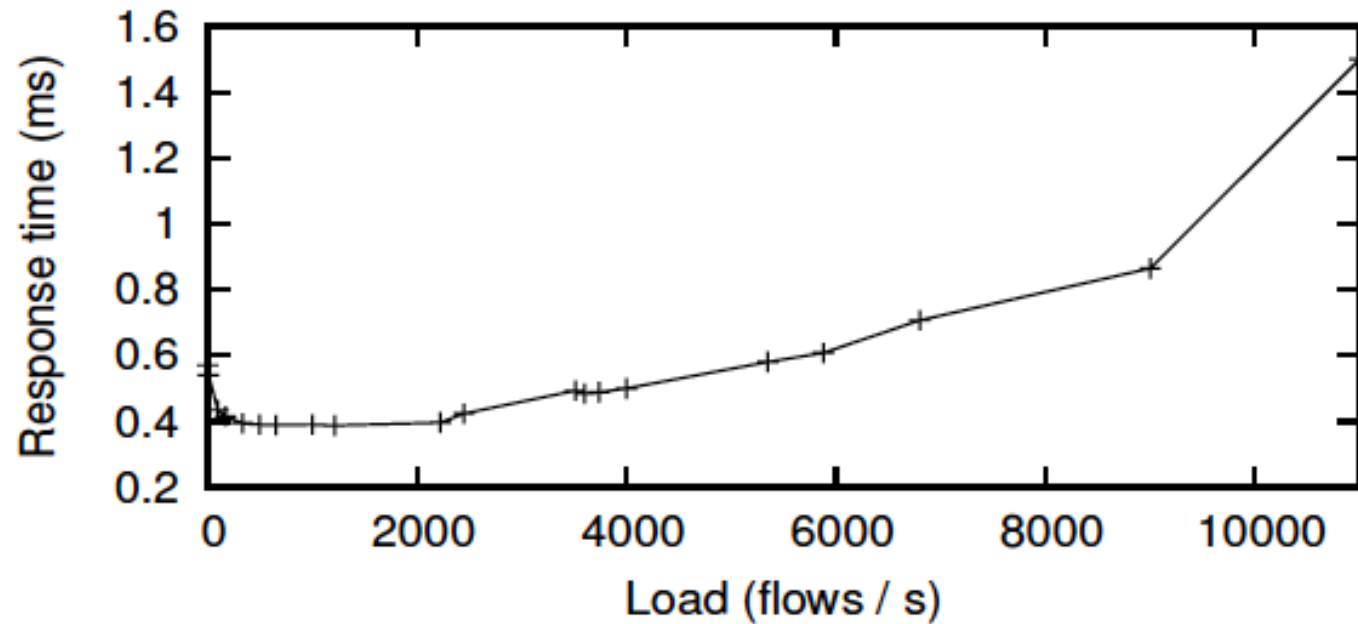


It is also worth noting that the flow table can be several orders-of-magnitude smaller than the forwarding table in an equivalent Ethernet switch. In an Ethernet switch, the table is sized to minimize broadcast traffic: as switches flood during learning, this can swamp links and makes the network less secure.<sup>5</sup> As a result, an Ethernet switch needs to remember all the addresses it's likely to encounter; even small wiring closet switches typically contain a million entries. Ethane Switches, on the other hand, can have much smaller

two-way hashing scheme [9]. A typical commercial enterprise Ethernet switch today holds 1 million Ethernet addresses (6MB, but larger if hashing is used), 1 million IP addresses (4MB of TCAM),

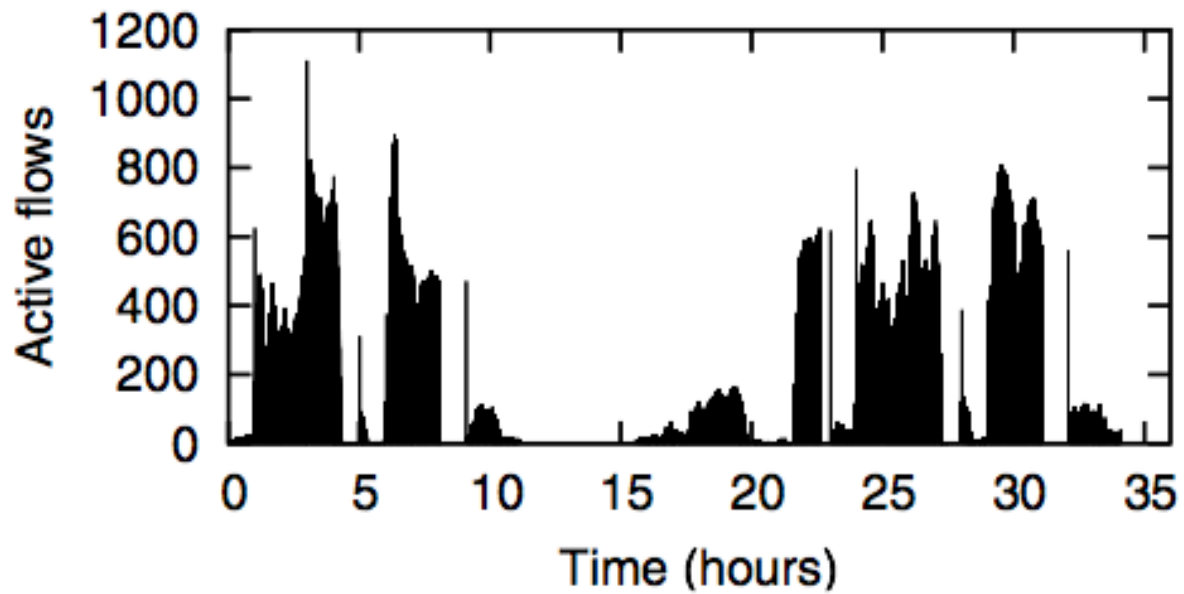
**Table 1.** Scalability Table

Name	WS-SUP720-3B	WS-SUP720-3BXL	VS-S720-10G-3C *	VS-S720-10G-3CXL*
MAC Entries	64,000	64,000	96,000	96,000
Routes	256,000 (IPv4); 128,000 (IPv6)	1,000,000 (IPv4); 500,000 (IPv6)	256,000 (IPv4); 128,000 (IPv6)	1,000,000 (IPv4); 500,000 (IPv6)

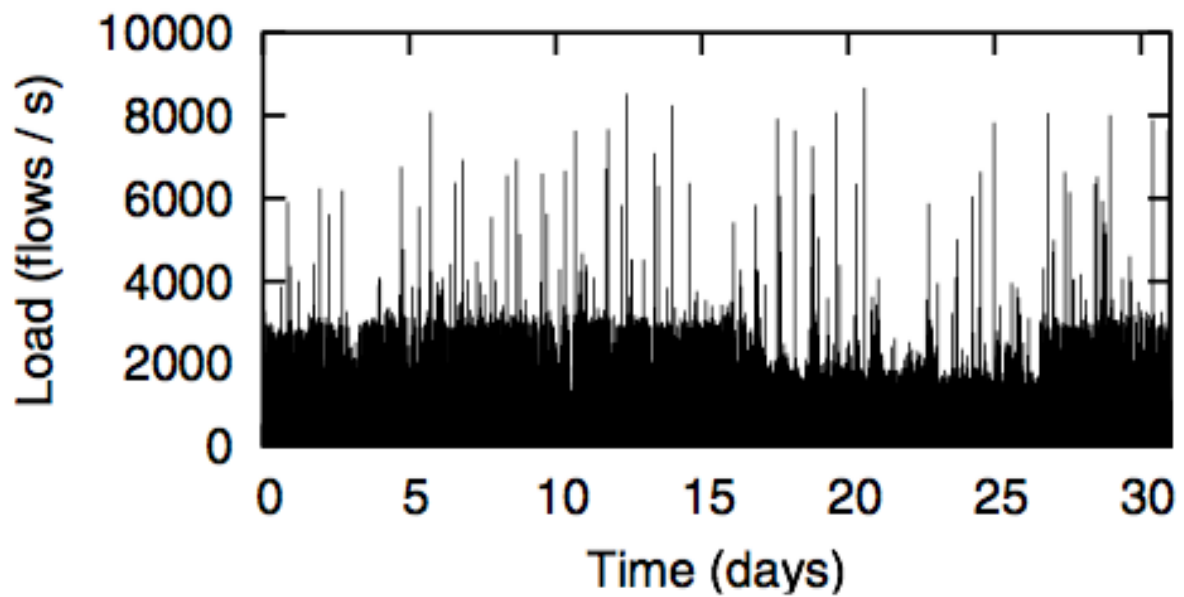


**Figure 6: Flow-setup times as a function of Controller load. Packet sizes were 64B, 128B and 256B, evenly distributed.**

heads. The Controller was configured with a policy file of 50 rules and 100 registered principles; routes were precalculated and cached. Under these conditions, the system could handle 650,845 bind events per second and 16,972,600 permission checks per second. The



**Figure 7: Active flows for LBL network [19].**



**Figure 8: Flow-request rate for Stanford network.**