



Towards Automated “Zero Day” Application Characterization: (Can we do better?)

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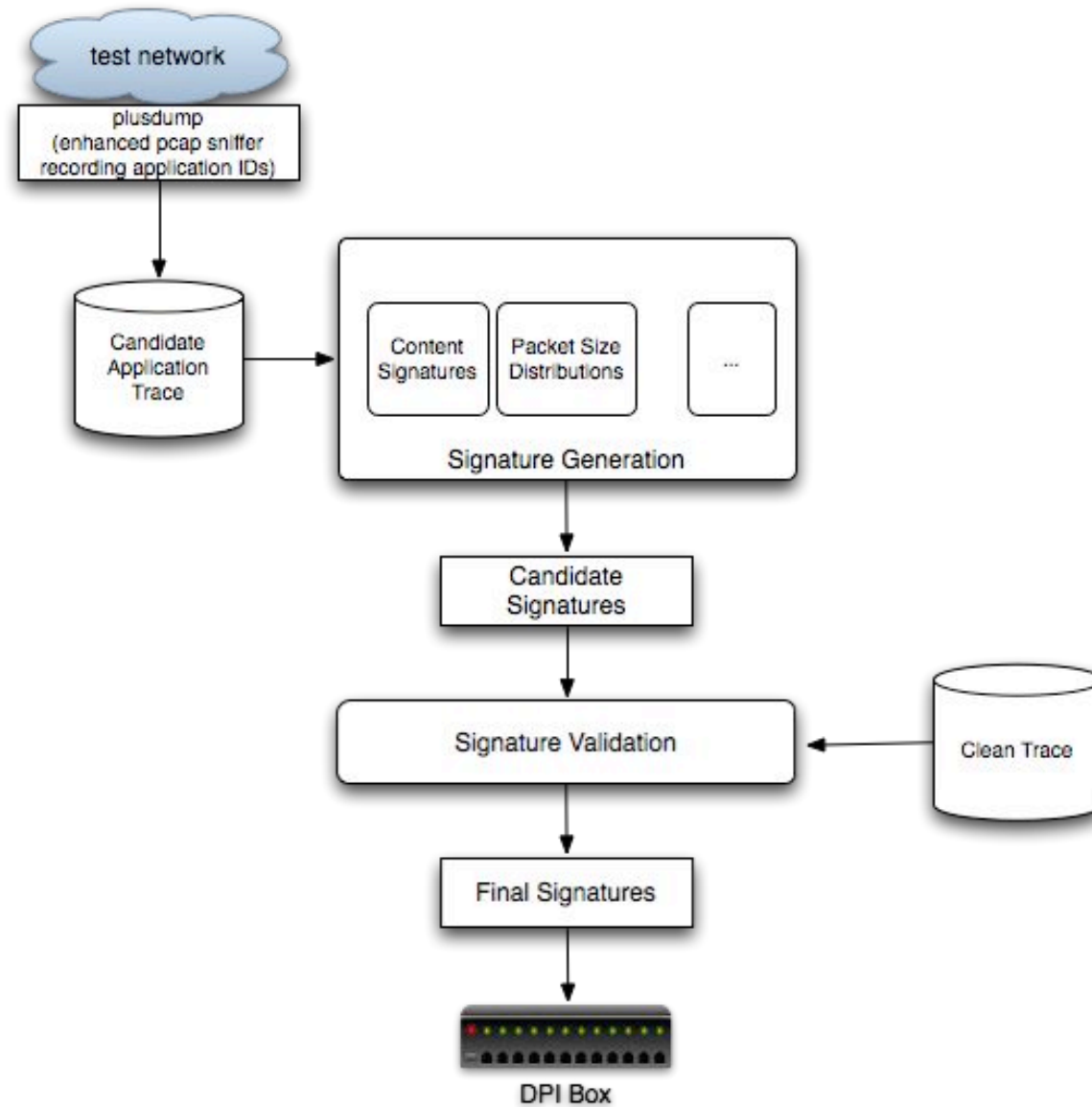
Jonathan M. Smith, *UPenn, USA*

Problem Statement & Goals



- **Pressing need for application characterization**
 - Maintain performance, control bandwidth usage
 - Protect against misbehaving or undesirable applications
 - Support networking research, or simply satisfy curiosity
- **Application characterization becoming complex**
 - New applications, application versions, new protocols
 - Obfuscation techniques: Skype, Joost, eMule, BitTorrent, etc.
- **Need for automation**
 - Tools to (semi-)automate characterization
 - Capability to rapidly develop+test characterization techniques
- **Our take on the problem**
 - Build upon techniques from zero-day attack sig generation
 - Strategic game between obscuring and revealing party

CUB4: Architecture Overview



Application content fingerprinting

- **Inspired by zero-day worm fingerprinting work**
 - [Akritidis2005,Earlybird2004, Autograph2005]
 - Use of rabin fingerprints over sliding window, with careful max-hit/max-size scoring threshold

ARBITRARY SENBELUTEK TENCE

Fingerprint = 11001001

ARANDOM BELUTEK STRINGABCD

Fingerprint = 11001001

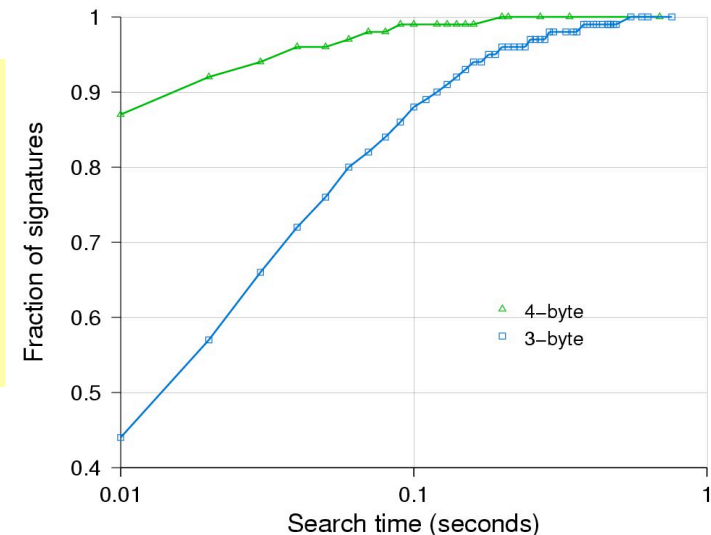
Signature Validation approach

- **Main idea: test candidate signatures against known traffic of other apps**
- **In a nutshell:**
 - Lots of data: for FPs in the order of 10^{-6} , need $\gg 10^6$ samples (1-100 GB)
 - Assume 1-100 signatures to test, need search time < 1 second
 - Approach: record packet trace, maintain carefully designed index for fast lookups
- **Trade-off:** need around 6x space to perform efficient lookups

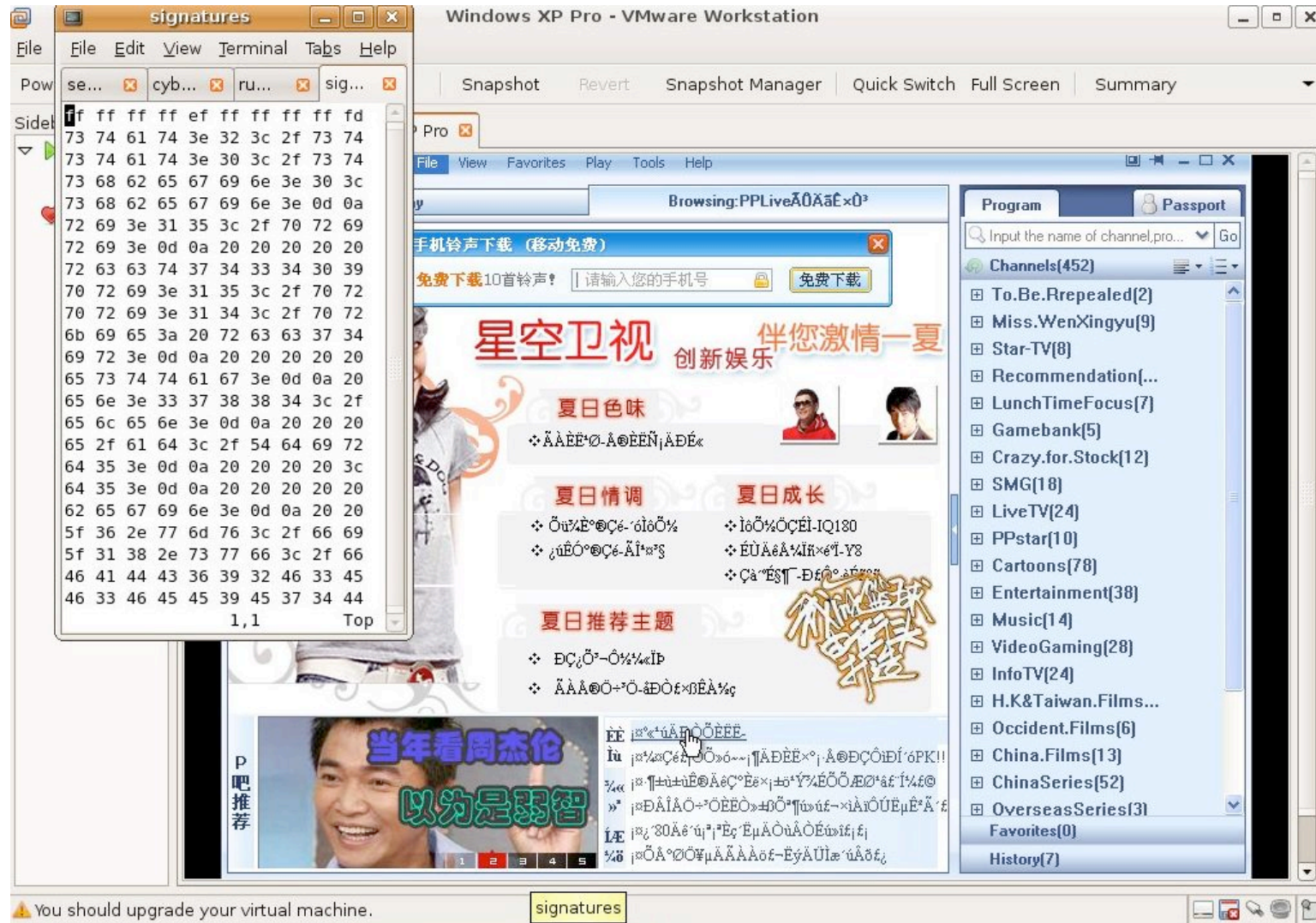
Search time: 88% < 0.01 sec, 99% < 0.1 sec, 99.99% < 1 sec,
known “bad” strings up to 2 seconds

Comparison: testing of PPLive signatures against 100GB
trace took 5 hours

Space: for N-byte elements, 2^{8N} bitmap + 6x trace size



Example signatures: PPLive



Application signature generation: can we go live?

•Reformulating the problem:

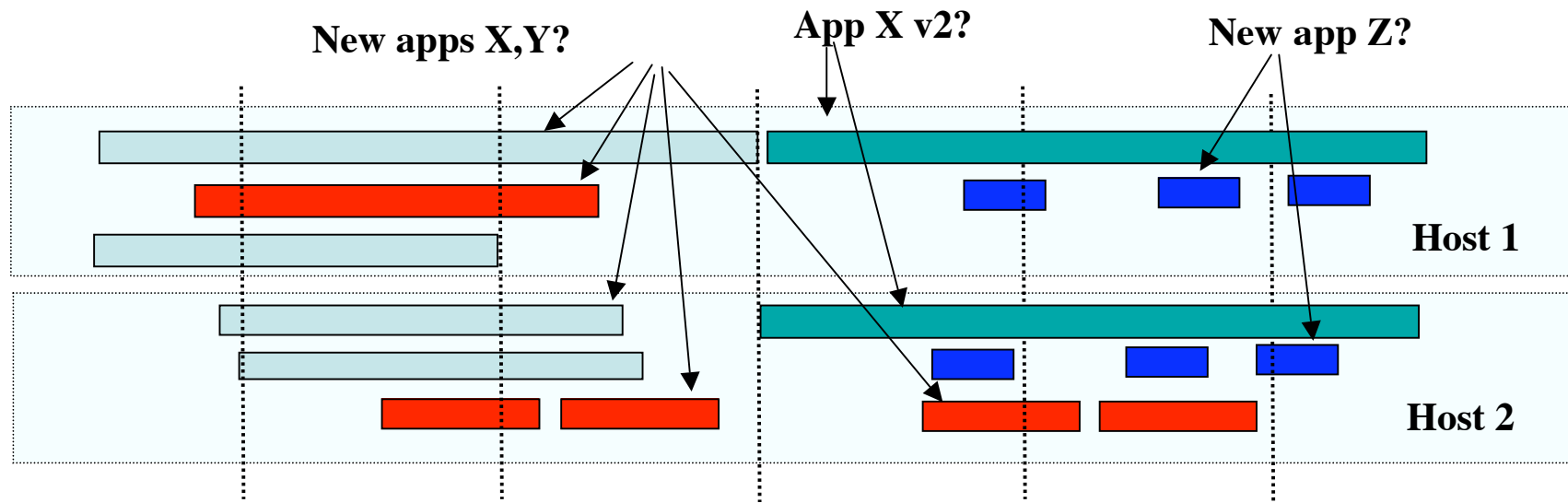
- So far, feeding our tool with “clean” offline traces
- Can we apply the same method to “live” traffic?



*“As we know, there are **known knowns**. There are things we know we know. We also know there are **known unknowns**. That is to say we know there are some things we do not know. But there are also **unknown unknowns**, the ones we don't know we don't know.”*

Donald Rumsfeld. Feb. 12, 2002,
Department of Defense news briefing

Application signature generation: can we go live?



- **Approach:**

- Split traces per host, and by time
- Determine joint fingerprints: flows of same host/ among hosts
- Compute likely set of applications based on the above data

- **High startup cost:**

- If new applications pop up one at a time, we may have a chance
- If we're looking at a link with 30-40% unclassified traffic from >20 applications, problem is somewhat more challenging

Beyond content: packet size

```
11:02:39.249981 IP 4.71.174.175.4166 > 58.185.58.82.49335: UDP, length 941
11:02:39.256917 IP 4.71.174.189.4166 > 58.185.58.82.49335: UDP, length 940
11:02:40.017697 IP 4.71.174.150.4166 > 58.185.58.82.49335: UDP, length 11
11:02:40.026990 IP 4.71.174.175.4166 > 58.185.58.82.49335: UDP, length 11
11:02:40.034641 IP 4.71.174.158.4166 > 58.185.58.82.49335: UDP, length 11
11:02:40.047234 IP 4.71.174.175.4166 > 58.185.58.82.49335: UDP, length 1057
11:02:40.291110 IP 4.71.174.158.4166 > 58.185.58.82.49335: UDP, length 1057
11:02:40.297065 IP 4.71.174.175.4166 > 58.185.58.82.49335: UDP, length 1058
11:02:40.312826 IP 4.71.174.150.4166 > 58.185.58.82.49335: UDP, length 1058
11:02:40.322903 IP 4.71.174.158.4166 > 58.185.58.82.49335: UDP, length 1058
```

Joost

```
06:48:13.610291 IP 131.111.218.93.64692 > 221.134.2.109.4662: S 687196497:687196497(0) >
06:48:14.536485 IP 221.134.2.109.4662 > 131.111.218.93.64692: S 1994439214:1994439214(0) >
06:48:14.539900 IP 131.111.218.93.64692 > 221.134.2.109.4662: P 1:121(120)
06:48:15.896222 IP 221.134.2.109.4662 > 131.111.218.93.64692: P 1:107(106)
06:48:16.117231 IP 131.111.218.93.64692 > 221.134.2.109.4662: P 121:230(109)
06:48:16.716196 IP 221.134.2.109.4662 > 131.111.218.93.64692: P 107:215(108)
06:48:16.776250 IP 131.111.218.93.64692 > 221.134.2.109.4662: P 230:271(41)
06:48:17.256531 IP 221.134.2.109.4662 > 131.111.218.93.64692: P 215:336(121)
06:48:17.339697 IP 131.111.218.93.64692 > 221.134.2.109.4662: P 271:293(22)
06:48:18.006485 IP 221.134.2.109.4662 > 131.111.218.93.64692: P 336:342(6)
06:48:18.038964 IP 131.111.218.93.64692 > 221.134.2.109.4662: P 293:309(46)
06:48:18.527099 IP 221.134.2.109.4662 > 131.111.218.93.64692: . 342:1802(1460)
06:48:18.596903 IP 221.134.2.109.4662 > 131.111.218.93.64692: P 1802:2942(1140)
```

**Obfuscated
eMule**

Behavioral analysis: packet size distributions

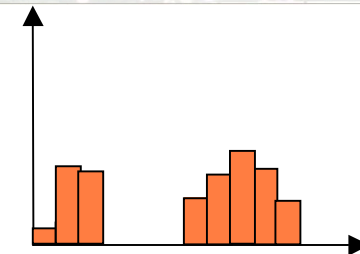
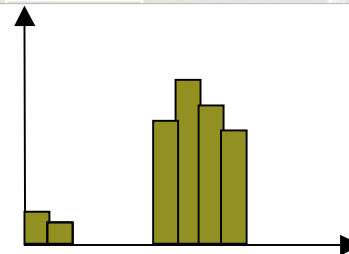
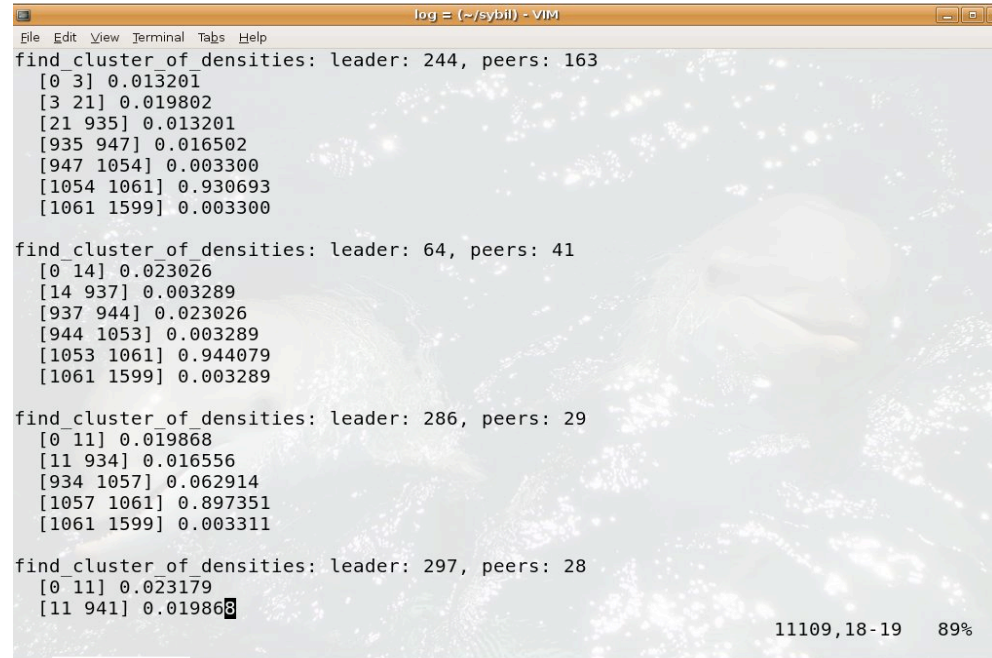
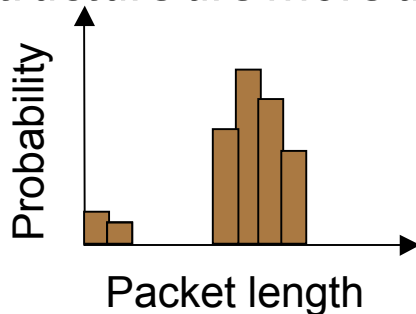
• Experiment: Packet size histogram over sliding window

- Various goodness-of-fit tests
- Fingerprinted Joost

Obtained signatures.

We can see a few representative probability density functions describing packet length.

- Other heuristics with more structure are more accurate



Increasing accuracy through taint propagation



- **Basic Idea:**

- If a flow $\langle \text{src}, \text{srcP}, \text{dstIP}, \text{dstP} \rangle$ previously identified to belong to app X, then any flow $\langle *, *, \text{dstIP}, \text{dstP} \rangle$ is very likely to belong to app X

- **Widely applicable:**

- Any application that operates over TCP and advertises “server”-side ports is relevant
- Performance gains likely to be higher at aggregation points
- Reminiscent of blacklisting, but this is not IP based: port # is the key!

- **Implications for heuristic design:**

- Tainting can compensate for low detection rates as long as the $\langle \text{dstIP}, \text{dstPort} \rangle$ pair handles multiple connections
- But need to be careful with false positive amplification
- On the other hand, we might tune down sensitivity to avoid FPs, as tainting will compensate for that in terms of detection rate

Some ongoing work



- **Incorporating structural features in packet size heuristics**
 - Much richer than simple packet size distribution analysis
 - Good preliminary results for obfuscated eMule and encrypted BitTorrent
- **Instrumenting endpoint software**
 - Current focus: use of Argos processor emulator, may need more lightweight
 - Looking for constants in memory that make it into packets
 - One step closer to fully automated fingerprinting
 - This would include cases where an application is updated, etc.
- **Detecting deep architectural properties of applications**
 - Adaptive applications through FEC-blowup or codec-switching
 - P2P through their incentive mechanisms
 - Much more labor-intensive to fingerprint, but also harder to circumvent
- **Active fingerprinting of likely “server” endpoints**
 - Even sending junk sometimes results in well-formed responses

Going adversarial



- **Two party model:**
 - Obscuring party (app dev) vs. revealing party: (researchers, DPI vendors)
- **This is already happening -- strategies we have seen:**
 - Avoid well known ports (too many to list)
 - Encrypt/obfuscate content (Skype, BitTorrent, eMule)
 - Anti-debugging techniques (Skype)
- **Other strategies we're likely to see:**
 - Content signatures → embedding of other application signatures
 - Packet size heuristics → proper padding
 - Binary instrumentation → more anti-debugging/anti-VM/anti-...
 - Active fingerprinting → only responding to conn requests from "trusted" sources
 - Connection graph/volume analysis → more cover traffic, dummy connections
 - ... at the limit, there's stego and Tor-like approaches

... but for the time being, security through obscurity might have value (?)



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