Personal Namespaces

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"I want a thousand guitars I want pounding drums I want a million different voices Speaking in tongues"

Naming Problems

Naming network resources and services is a big mess

The Solution

• Name everything "peachtree"



Naming Problems

Naming network resources and services is a big mess

- Problem #1: names are obtuse
- Problem #2: names are hard to share

http://www.flickr.com/photo_zoom.gne?id=1131208946&size=o&context=photostream

• Problem #3: names are gloablly unique, but ambiguous to people

- What is asu.edu?
 - Arkansas State University?
 - Alaska State University?
 - Arizona State University?
 - That fabulous little Division I-AA school in Boone, NC?

 Problem #4: names are intolerant of location change

mallman@cs.ohiou.edu ma137591@ohiou.edu



mallman@grc.nasa.gov

mallman@icir.org

mark.allman@case.edu

- Problem #5: naming is under nobody's control
 - Service providers play a part
 - E.g., "www.blogspot.com"
 - Content providers play a part
 - E.g., "MyGreatVacationPictures.html"
 - Consumers play a part
 - E.g., "Joe's Blog" in the bookmarks list

- But, we cope
 - Address books
 - Bookmarks
 - "Mail this web page"
 - Clicking URLs
 - Google
 - Social networking sites
- Is there a better way?

Related Work

- Much naming literature
 - How to name hosts
 - How to name services
 - How to name data
 - Personal or group naming realms

• All good stuff but issues remain

A Naming Layer

- Perhaps what we need is a new over-arching namespace
 - Just an abstraction to existing namespaces

A Naming Layer (cont.)

- Give users' a way to name their own resources
 - Independent of resource/service location
 - With context sensitive names
 - Scoping defined by the user
 - Public vs. private



- Every pnames user gets a namespace
 - Identified by a namespace ID (NID)
 - NID is a hash of the public half of a locally generated keypair

Overview (cont.)

- Each namespace can contain:
 - Simple names
 - E.g., "calendar = webcal://cal.mallman...."
 - E.g., "email = mallman@icir.org"
 - E.g., "aim = myAlMhandle"
 - Pointers to other namespaces
 - E.g., "Joe = NID:7a6b623df1"

Example



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Example (cont.)

- Mark can use:
 - Dad:blog
- Wes can use:
 - Mark:vacation-pix
 - Mark:web

Sharing Names

- Scheme #1: sharing through a DHT
 - Hash key composed of NID, name and type
 - Value is the actual resource name

Sharing Names (cont.)

- This scheme works on a *per-name* basis, not a per-namespace basis
 - I.e., lookups do not retrieve all the names someone makes available

Sharing Names (cont.)

- Scheme #2: names could be swapped without using a DHT
 - However users now informally share data

Bootstrapping

- Key problem: NIDs are *more obscure* than any other sort of name we already use
 - Makes them *harder* to share
 - Makes pointer records crucial to the system

Bootstrapping (cont.)

- Swap NIDs any way information is shared now
 - Email signatures, web pages, vCards, etc.

- Could also setup *pnames* registeries where users could add pointers via some web page
 - E.g., "WellKnownRegistry:MarkAllman"

Reliability

• Adding a new naming layer adds a new point of failure

- We combat this in two ways:
 - DHTs are robust
 - Heavily use caching and pre-fetching

Security

- Names can be validated as belonging to the given namespace by checking signatures
- Given crypto usage is required we can encrypt records we want to share only privately
- DHT is robust to some kinds of DoS attacks since it is a *distributed* data structure
 - Of course, different attacks may leverage homogeneous software in a tightly coupled system

Final Thoughts

- Pnames is a simple new abstraction
 - Do we need it?
 - Is it solving a problem we have?
 - Are our current coping mechansisms Good Enough?

"Is there anybody alive out there?"

Extra Slides

A Different Namespace

- Pnames is not strictly scoped to people
 - E.g., organizations could build a namespace

- What are the implications?
- Who is authoritative if "google.com" resolves differently in different environments?

OpenDHT Performance



Brian Beck, Benjamin Golub, Scott Reid, Mark Allman. On the Performance of OpenDHT. October 2007. Under submission.

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