

# Malware: Viruses

***CS 161: Computer Security***

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# Inside a Modern HIDS (“AV”)

- URL/Web access blocking:
  - Prevent users from going to **known bad locations**
- Protocol scanning of network traffic (esp. HTTP)
  - Detect & block known **attacks**
  - Detect & block known **malware communication**
- Payload scanning
  - Detect & block known **malware**
- (Auto-update of signatures for these)
- **Cloud queries** regarding reputation
  - Who else has run this executable and with what results?
  - What’s known about the remote host / domain / URL?

# Inside a Modern HIDS, con't

- *Sandbox execution*
  - Run selected executables in constrained/monitored environment
  - Analyze:
    - System calls
    - Changes to files / registry
    - Self-modifying code (*polymorphism/metamorphism*)
- File scanning
  - Look for known malware that installs itself on disk
- Memory scanning
  - Look for known malware that **never appears on disk**
- Runtime analysis
  - Apply heuristics/signatures to execution behavior

# Inside a Modern NIDS

- Deployment **inside** network as well as at border
  - Greater visibility, including **tracking of user identity**
- Full **protocol analysis**
  - Including extraction of complex embedded objects
  - In some systems, 100s of known protocols
- Signature analysis (also behavioral)
  - **Known** attacks/vulnerabilities, malware communication, blacklisted hosts/domains
  - **Known** malicious payloads
  - Sequences/patterns of activity
- **Shadow execution** (e.g., Flash, PDF programs)
- Extensive logging (in support of **forensics**)
- Auto-update of signatures, blacklists; cloud queries

# Malware

# The Problem of Malware

- **Malware** = malicious code that runs on a victim's system
- How does it manage to run?
  - Attacks a network-accessible **vulnerable service**
  - **Vulnerable client** connects to remote system that sends over an attack (a **driveby**)
  - **Social engineering**: trick user into running/installing
  - “Autorun” functionality (esp. from plugging in USB device)
  - Slipped into a system component (at manufacture; compromise of software provider; substituted via **MITM**)
  - **Attacker with local access** downloads/runs it directly
    - Might include using a local “privilege escalation” exploit

# What Can Malware Do?

- Pretty much *anything*
  - Payload generally *decoupled* from how manages to run
  - Only subject to *permissions* under which it runs
- Examples:
  - Brag or exhort or extort (pop up a message/display)
  - Trash files (just to be nasty)
  - Damage hardware (!)
  - Launch external activity (spam, *click fraud*, DoS; *banking*)
  - Steal information (*exfiltrate*)
  - Keylogging; screen / audio / camera capture
  - Encrypt files (*ransomware*)
- Possibly delayed until condition occurs
  - “*time bomb*” / “*logic bomb*”

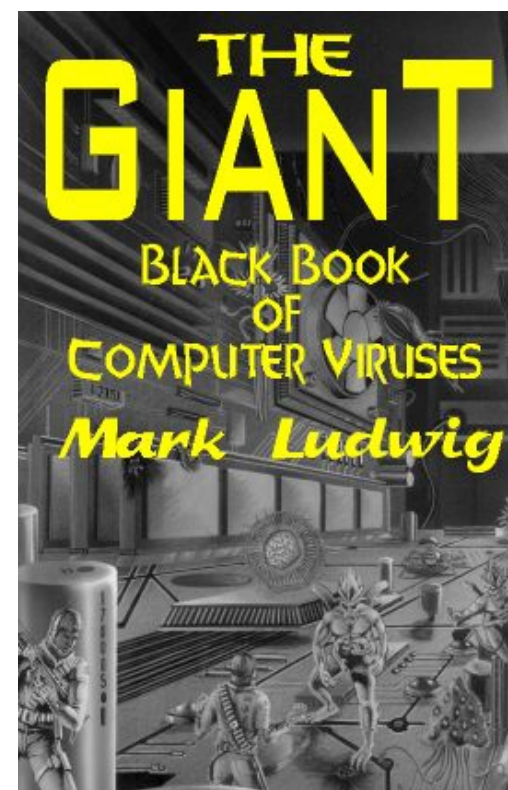
# Malware That Automatically Propagates

- **Virus** = code that **propagates** (replicates) across systems by arranging to have itself *eventually executed*, creating a **new additional instance**
  - Generally infects by altering **stored** code
- **Worm** = code that **self-propagates**/replicates across systems by arranging to have itself *immediately executed* (creating **new addl. instance**)
  - Generally infects by altering **running** code
  - No user intervention required
- (Note: line between these isn't always so crisp; plus some malware incorporates both approaches)



# The Problem of Viruses

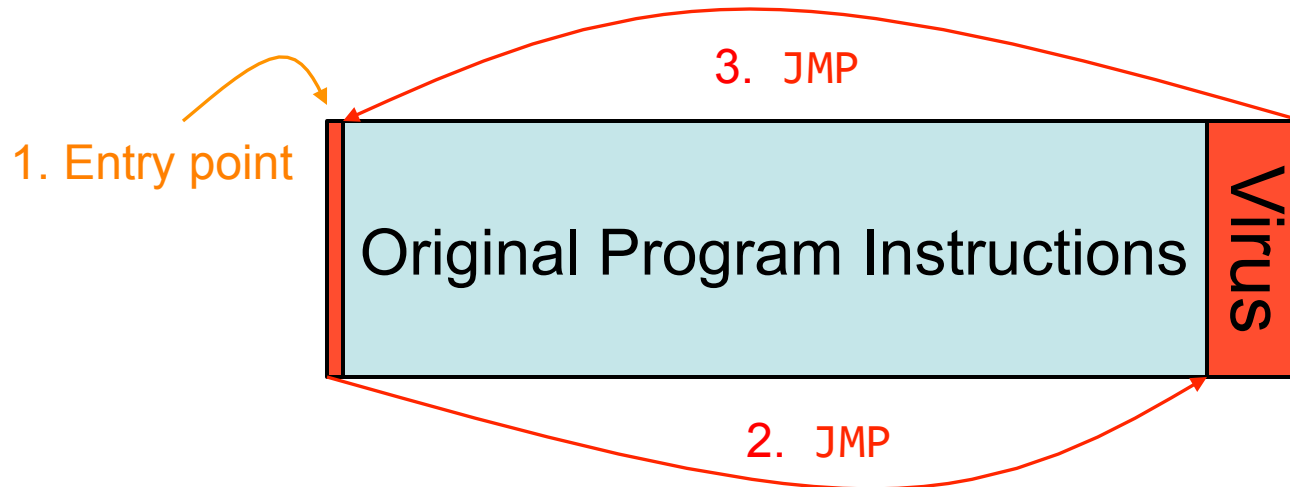
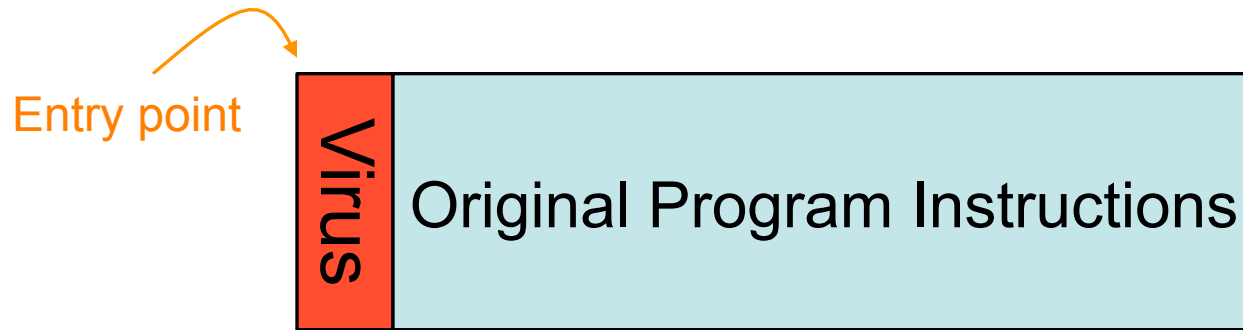
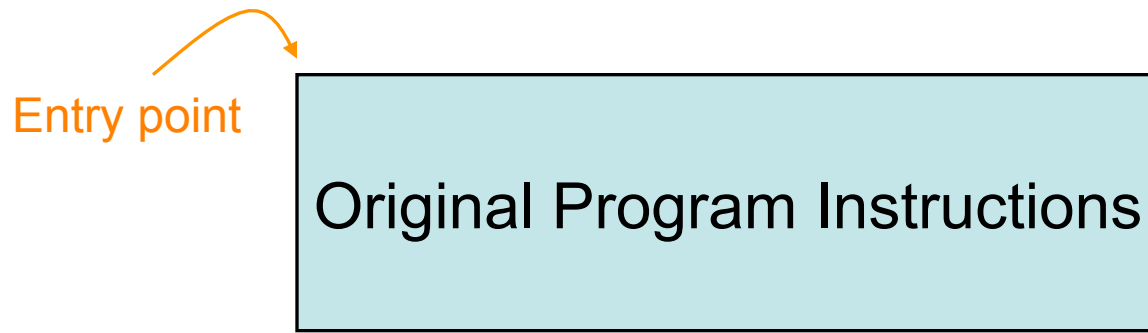
- Opportunistic = code will **eventually** execute
  - Generally due to **user action**
    - Running an app, booting their system, opening an attachment
- Separate notions: how it **propagates** vs. what else it does when executed (**payload**)
- General infection strategy: find some code lying around, **alter it** to include the virus
- Have been around for **decades** ...
  - ... resulting **arms race** has heavily influenced evolution of modern malware



# Propagation

- When virus runs, it looks for an **opportunity** to infect additional systems
- One approach: look for USB-attached thumb drive, alter any executables it holds to include the virus
  - Strategy: when drive later attached to **another** system & altered executable runs, it locates and infects executables on **new** system's hard drive
- **Or:** when user sends email w/ attachment, virus **alters attachment** to add a copy of itself
  - Works for attachment types that include **programmability**
  - E.g., Word documents (macros)
  - Virus can also send out such email proactively, using user's address book + enticing subject (“**I Love You**”)

*autorun is handy here!*



Original program instructions can be:

- Application the user runs
- Run-time library / routines resident in memory
- Disk blocks used to boot OS
- Autorun file on USB device
- ...

Other variants are possible; whatever manages to get the virus code executed

# Detecting Viruses

- Signature-based detection
  - Look for bytes corresponding to injected virus code
  - High utility due to **replicating nature**
    - If you capture a virus V on one system, by its nature the virus will be trying to infect *many other systems*
    - Can protect those other systems by installing recognizer for V
- Drove development of **multi-billion \$\$ AV industry** (AV = “antivirus”)
  - So many **endemic** viruses that detecting well-known ones becomes a “*checklist item*” for security audits
- Using signature-based detection also has de facto utility for (glib) **marketing**
  - Companies compete on number of signatures ...
    - ... rather than their quality (harder for customer to assess)

SHA256: 58860062c9844377987d22826eb17d9130dceaa7f0fa68ec9d44dfa435d6ded4

File name: cc8caa3d2996bf0360981781869f0c82.exe

Detection ratio: 11 / 62

Analysis date: 2017-04-18 22:28:27 UTC ( 56 minutes ago )



Analysis

File detail

Relationships

Additional information

Comments 4

Votes

Behavioural information

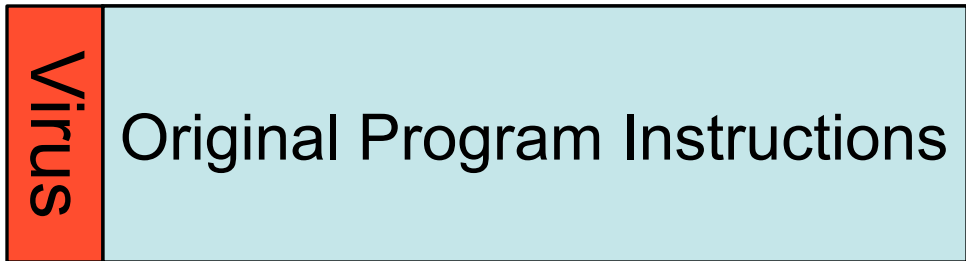
Antivirus	Result	Update
Avira (no cloud)	TR/Crypt.ZPACK.atbin	20170418
CrowdStrike Falcon (ML)	malicious_confidence_100% (W)	20170130
DrWeb	Trojan.PWS.Panda.11620	20170418
Endgame	malicious (moderate confidence)	20170413
ESET-NOD32	a variant of Win32/GenKryptik.ACKE	20170418
Invincea	virus.win32.ramnit.ah	20170413
Kaspersky	Trojan.Win32.Yakes.tavs	20170418
Palo Alto Networks (Known Signatures)	generic.ml	20170418
TrendMicro-HouseCall	Suspicious_GEN.F47V0418	20170418
Webroot	W32.Malware.Gen	20170418
ZoneAlarm by Check Point	Trojan.Win32.Yakes.tavs	20170418
Ad-Aware	✓	20170418
AegisLab	✓	20170418

# Virus Writer / AV Arms Race

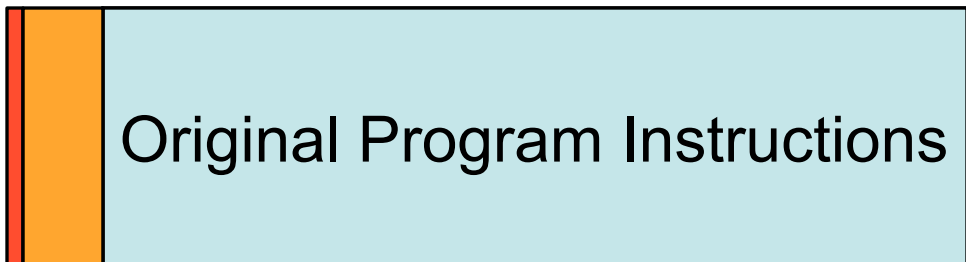
- If you are a virus writer and your beautiful new creations don't get very far because each time you write one, the AV companies quickly push out a signature for it ....
  - .... *What are you going to do?*
- Need to keep **changing** your viruses ...
  - ... or at least changing their appearance!
- How can you **mechanize** the creation of new instances of your viruses ...
  - ... so that whenever your virus propagates, what it injects as a copy of itself **looks different?**

# Polymorphic Code

- We've already seen technology for creating a representation of data apparently completely unrelated to the original: **encryption!**
- Idea: every time your virus propagates, it inserts a **newly encrypted copy** of itself
  - Clearly, encryption needs to vary
    - Either by using a different key each time
    - Or by including some random initial padding (like an IV)
  - Note: weak (but simple/fast) crypto algorithm works fine
    - No need for truly strong encryption, just **obfuscation**
- When injected code runs, it decrypts itself to obtain the original functionality



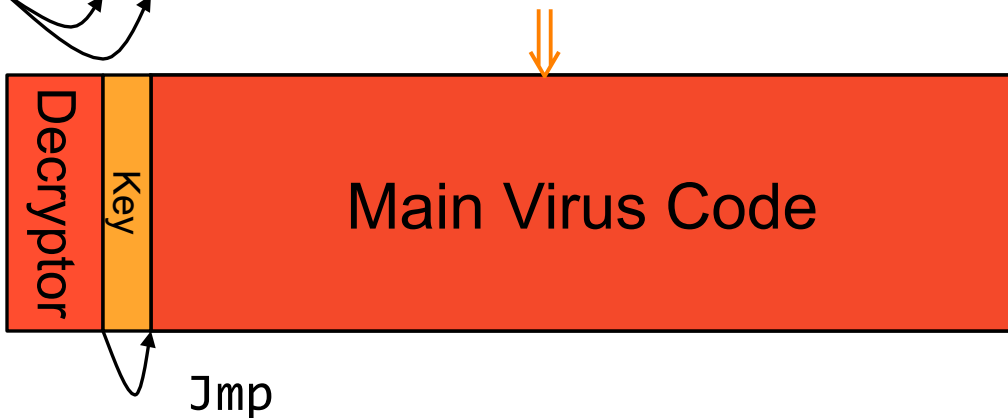
Instead of this ...



Virus has *this* **initial** structure



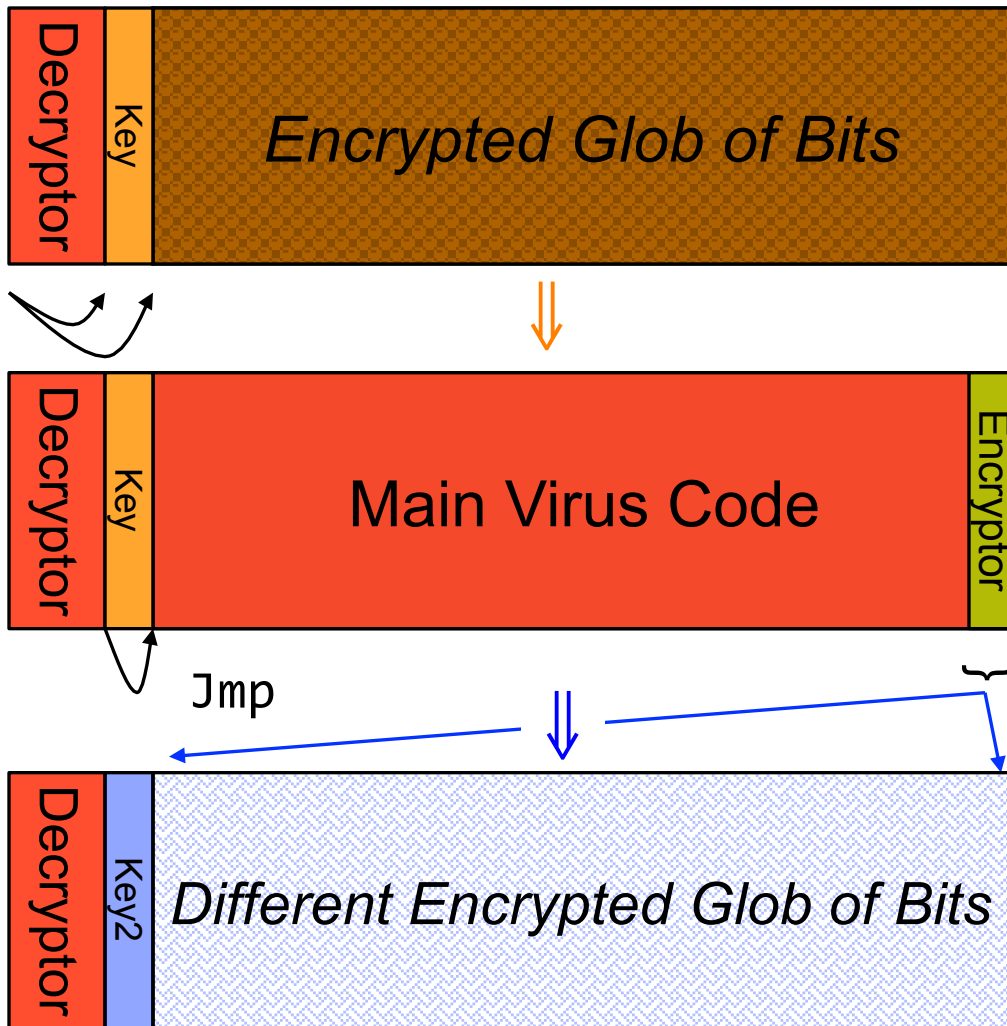
When executed, decryptor applies key to decrypt the glob ...



... and jumps to the decrypted code once stored in memory



# Polymorphic Propagation



Once running, virus uses an *encryptor* with a **new key** to propagate

New virus instance bears **little resemblance** to original

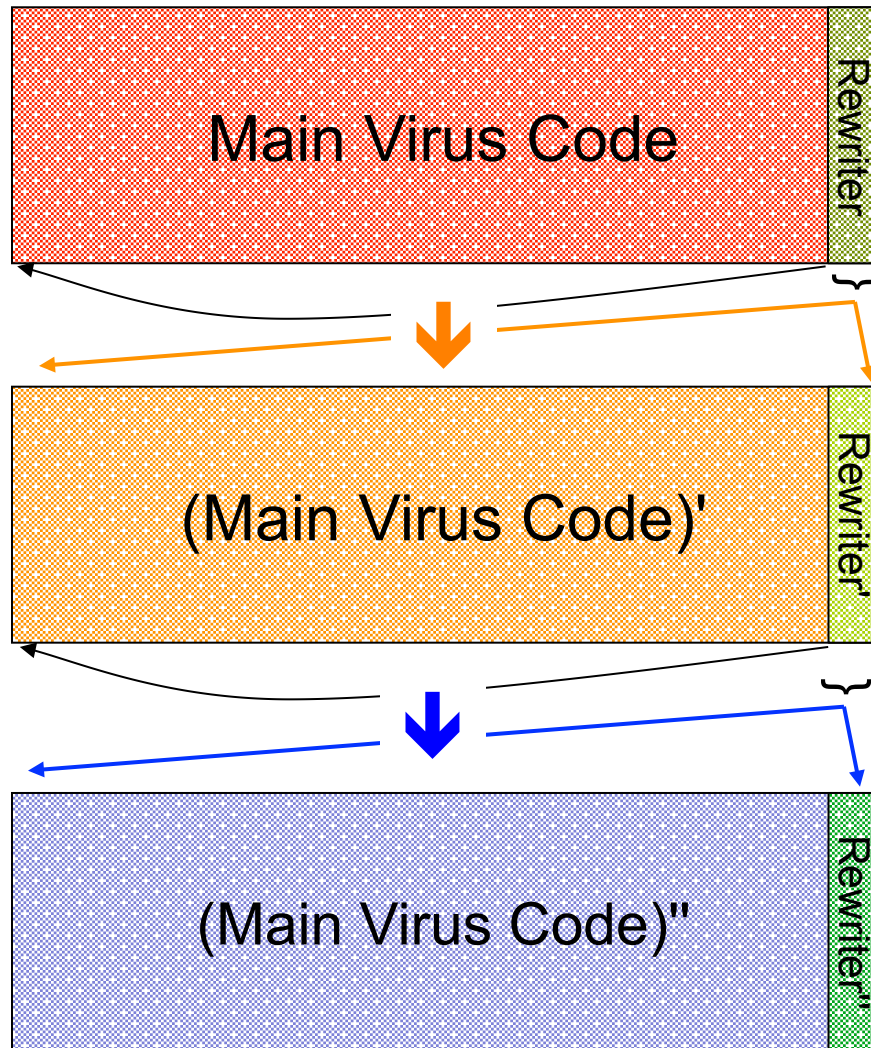
# Arms Race: Polymorphic Code

- Given polymorphism, how might we then detect viruses?
- Idea #1: use narrow sig. that targets decryptor
  - Issues?
    - Less code to match against  $\Rightarrow$  more **false positives**
    - Virus writer spreads decryptor across existing code
- Idea #2: execute (or statically analyze) suspect code to see if it decrypts!
  - Issues?
    - Legitimate “*packers*” perform similar operations (decompression)
    - How long do you let the new code execute?
      - If decryptor only acts after lengthy legit execution, difficult to spot
- Virus-writer countermeasures?

# Metamorphic Code

- Idea: every time the virus propagates, generate *semantically different* version of it!
  - Different semantics only at immediate level of execution; higher-level semantics remain same
- How could you do this?
- Include with the virus a **code rewriter**:
  - Inspects its own code, generates random variant, e.g.:
    - Renumber registers
    - Change order of conditional code
    - Reorder operations not dependent on one another
    - Replace one low-level algorithm with another
    - Remove some do-nothing **padding** and replace with different do-nothing padding (“chaff”)
      - Can be very complex, legit code ... if it’s never called!

# Metamorphic Propagation



When ready to propagate, virus invokes a randomized *rewriter* to construct **new but semantically equivalent** code (including the rewriter)

# Detecting Metamorphic Viruses?

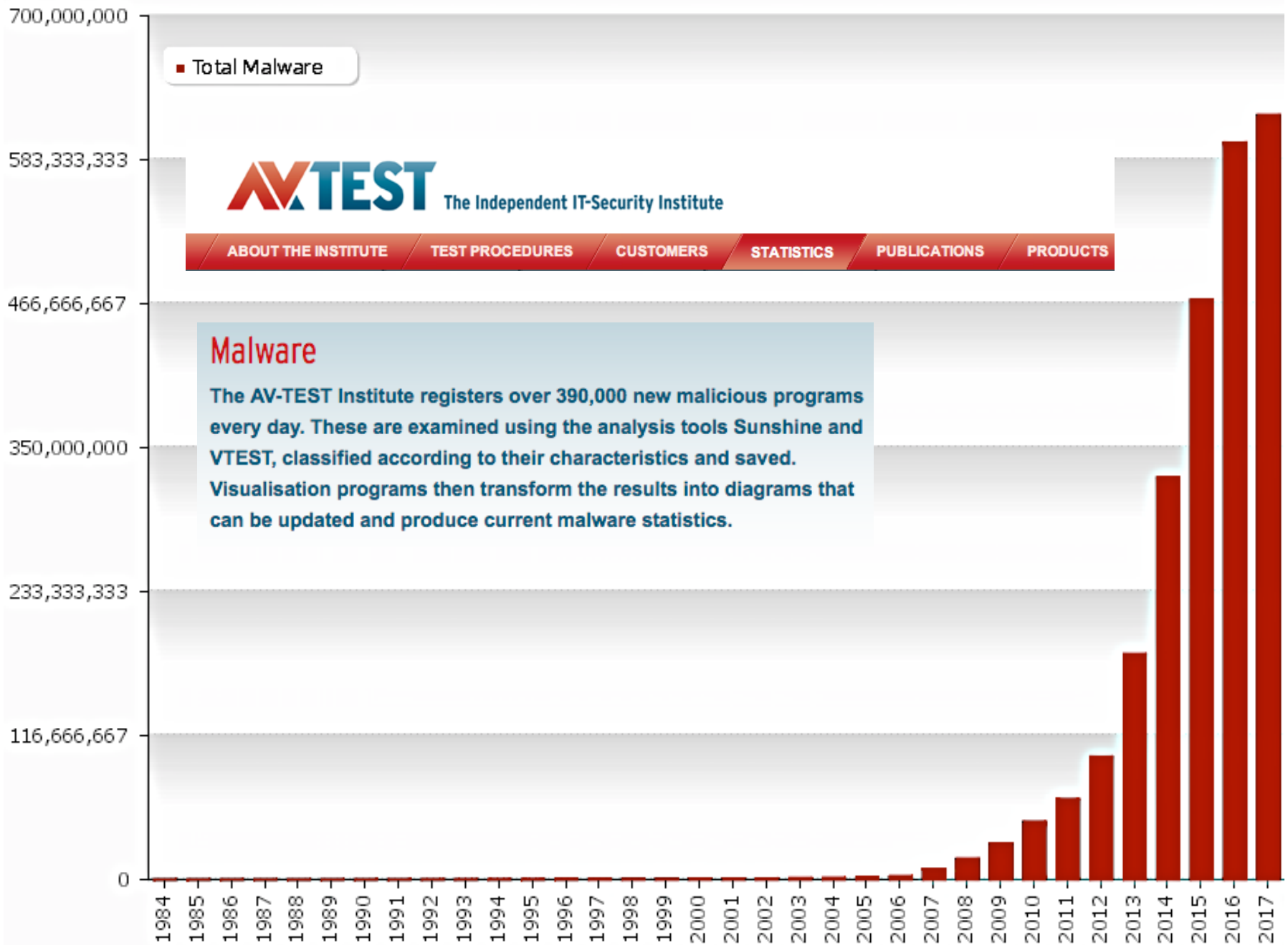
- Need to analyze execution **behavior**
  - Shift from **syntax** (*appearance* of instructions) to **semantics** (*effect* of instructions)
- Two stages: (1) AV company analyzes new virus to find **behavioral signature**; (2) AV software on end systems analyze suspect code to test for match to signature
- What countermeasures will the virus writer take?
  - **Delay analysis** by taking a long time to manifest behavior
    - Long time = await particular condition, or even simply clock time
  - Detect that execution occurs in an **analyzed environment** and if so behave differently
    - E.g., test whether running inside a debugger, or in a Virtual Machine
- Counter-countermeasure?
  - AV analysis looks for these tactics and skips over them
- Note: attacker has edge as ***AV products supply an oracle***

**5 Minute Break**

**Questions Before We Proceed?**

# How Much Malware Is Out There?

- A final consideration re polymorphism and metamorphism:
  - Presence can lead to **mis-counting** a single virus outbreak as instead reflecting 1,000s of *seemingly different* viruses
- Thus **take care** in interpreting vendor **statistics** on malware varieties
  - (Also note: public perception that huge malware populations exist is *in the vendors' own interest*)





# Infection Cleanup

- Once malware detected on a system, how do we get **rid** of it?
- May require restoring/repairing many files
  - This is part of what AV companies sell: per-specimen disinfection procedures
- What about if malware executed with **administrator privileges**?
  - *“nuke the entire site from orbit. It’s the only way to be sure”*  
- ALIENS
  - i.e., **rebuild** system from **original media + data backups**
- Malware may include a **rootkit**: *kernel patches* to **hide its presence** (its existence on disk, processes)

# Infection Cleanup, con't

- If we have complete source code for system, we could rebuild from that instead, couldn't we?
- No!
- Suppose forensic analysis shows that virus introduced a **backdoor** in `/bin/login` executable
  - (Note: this threat isn't specific to viruses; applies to any malware)
- Cleanup procedure: rebuild `/bin/login` from source ...



/bin/login  
source code

Compiler

Regular compilation  
process of building login  
binary from source code

/bin/login  
executable

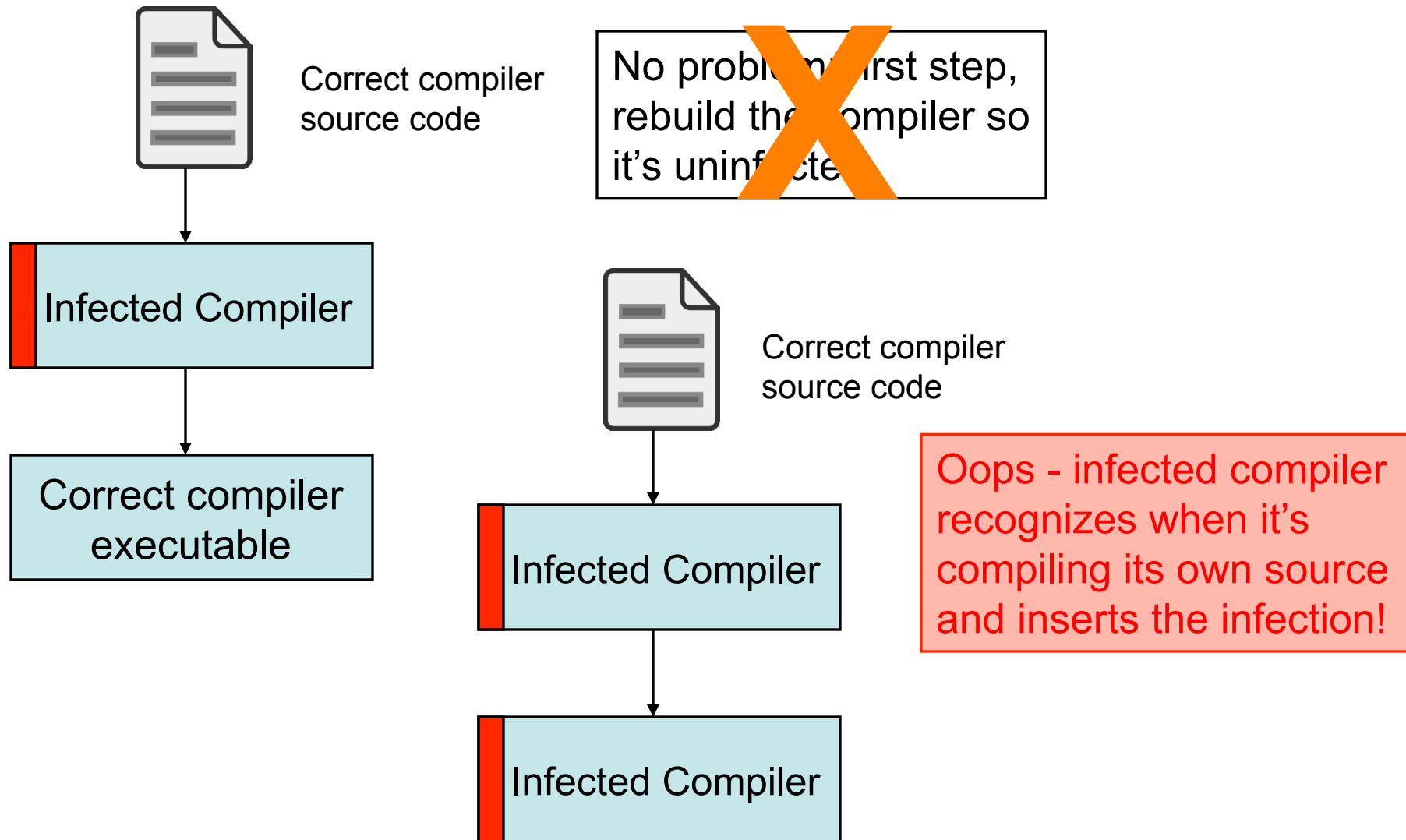


/bin/login  
source code

Compiler

**Infected compiler**  
*recognizes when it's  
compiling /bin/login  
source and inserts extra  
back door when seen*

/bin/login  
executable



**No** amount of careful source-code scrutiny can prevent this problem. And if the *hardware* has a back door ...

*Reflections on Trusting Trust*  
Turing-Award Lecture, Ken Thompson, 1983