#### **Network Control**

## CS 161: Computer Security Prof. Vern Paxson

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## **Review: Sniffing & Spoofing**







Because of this, B can spoof DHCP offers to A, and vice versa. But no one else can, because the requests stay within A's subnet.





... other than for local traffic such as between A & B







#### **Network Diagram**





The use of randomized fields in TCP & DNS make this very hard.

## Controlling Networks Using Firewalls

#### **Controlling Networks ... On The Cheap**

- Motivation: How do you harden a set of systems against external attack?
  - Key Observation:
    - The more network services your machines run, the greater the risk
  - Due to larger attack surface
- One approach: on each system, turn off unnecessary network services
  - But you have to know all the services that are running
  - And sometimes some trusted remote users still require access
- Plus key question of scaling
  - What happens when you have to secure 100s/1000s of systems?
  - Which may have different OSs, hardware & users ...
  - Which may in fact not all even be identified ...

# **Taming Management Complexity**

- Possibly more scalable defense: Reduce risk by blocking *in the network* outsiders from having unwanted access your network services
  - Interpose a firewall that traffic to/from the outside must traverse
  - Chokepoint can cover 1000s of hosts



# **Selecting a Security Policy**

- Effectiveness of firewall relies on deciding what policy it should implement:
  - Who is allowed to talk to whom, accessing what service?
- Distinguish between inbound & outbound connections
  - Inbound: attempts by external users to connect to services on internal machines
  - Outbound: internal users to external services
  - Why? Because fits with a common *threat model*

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  - Outbound: internal users to external services
  - Why? Because fits with a common *threat model*
- Conceptually simple *access control policy*:
  - Permit inside users to connect to any service
  - External users restricted:
    - Permit connections to services meant to be externally visible
    - Deny connections to services not meant for external access

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#### How To Treat Traffic Not Mentioned in Policy?

• **Default Allow**: start off permitting external access to services

Shut them off as problems recognized

- **Default Deny:** ✓ off permitting just a few known, well-secured services
  - Add more when users complain (and mgt. approves)
- Pros & Cons?

In general, use Default Deny

- Flexibility vs. conservative design
- Flaws in Default Deny get noticed more quickly / less painfully

## **Packet Filters**

- Most basic kind of firewall is a packet filter
  - Router with list of access control rules
  - Router checks each received packet against security rules to decide to forward or drop it
  - Each rule specifies which packets it applies to based on a packet's header fields (stateless)
    - Specify source and destination IP addresses, port numbers, and protocol names, or wild cards



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    - Specify source and destination IP addresses, port numbers, and protocol names, or wild cards
    - Each rule specifies the *action* for matching packets: ALLOW or DROP (aka DENY)

<ACTION> <PROTO> <SRC:PORT> -> <DST:PORT>

– First listed rule has *precedence* 

### **Examples of Packet Filter Rules**

allow tcp 4.5.5.4:1025 -> 3.1.1.2:80

- States that the firewall should permit any TCP packet that's:
  - from Internet address 4.5.5.4 and
  - using a source port of 1025 and
  - destined to port 80 of Internet address 3.1.1.2

deny tcp 4.5.5.4:\* -> 3.1.1.2:80

• States that the firewall should **drop** any TCP packet like the above, regardless of source port

#### **Examples of Packet Filter Rules**

deny tcp 4.5.5.4:\* -> 3.1.1.2:80 allow tcp 4.5.5.4:1025 -> 3.1.1.2:80

 In this order, the rules won't allow any TCP packets from 4.5.5.4 to port 80 of 3.1.1.2

allow tcp 4.5.5.4:1025 -> 3.1.1.2:80 deny tcp 4.5.5.4:\* -> 3.1.1.2:80

 In this order, the rules allow TCP packets from 4.5.5.4 to port 80 of 3.1.1.2 only if they come from source port 1025

#### **Firewall Considerations**

- Firewalls can have 1000s of filtering rules like these
  Easy to introduce subtle errors 😕
- Provide not only security but also policy enforcement
  E.g. do not allow company systems to access file-sharing sites
- Modern firewalls operate in a stateful fashion
  - Make Yes/No decisions upon establishment of a connection/flow
    - For Yes decisions, add 4-tuple to a *connection table* consulted for future traffic
    - Drop arriving non-establishment packet if not in table
- An important example of a *reference monitor*



#### **Security Principle:** *Reference Monitors*

- Firewalls embody useful principles that are applicable elsewhere in computer security
  - Optimized for enforcing particular kind of access control policy
  - Chokepoint notion makes enforcement possible
- A reference monitor examines every request to access a controlled resource (an object) and determines whether to allow request



### **Reference Monitor Security Properties**

- Always invoked
  - Complete mediation property: all security-relevant operations must be mediated by RM
  - RM should be invoked on every operation controlled by access control policy
- Tamper-resistant
  - Maintain RM integrity (no code/state tampering)
- Verifiable
  - Can verify RM operation (correctly enforces desired access control policy)
    - Requires extremely **simple** RM
    - We find we can't verify correctness for systems with any appreciable degree of complexity

# Considering Firewalls as Reference Monitors

- Always invoked?
  - Place Packet Filter as an *in-path* element on chokepoint link for all internal-external communications
  - Packets only forwarded across link if firewall explicitly decides to do so after inspection

# **Potential Problems?**

- What if a user hooks up an unsecured wireless access point to their internal machine?
- Anyone who drives by with wireless-enabled laptop can gain access to internal network
  – Bypasses packet filter!
- Or: what if user brings an infected device onto the premises?
- To use a firewall safely, must ensure we've covered all links between internal and external/ untrusted networks with firewalls
  - Set of links known as the security perimeter

# **RM Property:** *Tamper-Resistant*

- Will this hold?
- Do not allow management access to firewall other than from specific hosts

– I.e., firewall itself needs firewalling

- Protect firewall's physical security
- Must also secure storage & propagation of configuration data

# **RM Property: Verifiable**

- Will this hold?
- Current practice:
  - Packet filter software too complex for feasible systematic verification ...
  - … and rulesets with 1,000s (!) of rules
- Result:
  - Bugs that allowed attackers to defeat intended security policy by sending unexpected packets that packet filter doesn't handle as desired
- In addition: challenging to ensure network topology does not allow internal access by untrusted devices

# Why Have Firewalls Been Successful?

- Central control easy administration and update
  - Single point of control: update one config to change security policies
  - Potentially allows rapid response
- Easy to deploy transparent to end users
  - Easy incremental/total deployment to protect 1,000's
- Addresses an important problem
  - Security vulnerabilities in network services are rampant
  - Easier to use firewall than to directly secure code ...

# **Firewall Disadvantages?**

- Functionality loss less connectivity, less risk
  - May reduce network's usefulness
  - Some applications don't work with firewalls
    - Two peer-to-peer users behind different firewalls
- The malicious insider problem
  - Deployment assumes insiders are trusted
    - Malicious insider (or *anyone gaining control of internal machine*) can wreak havoc
- Firewalls establish a *security perimeter* 
  - Like Eskimo Pies: "hard crunchy exterior, soft creamy center"
  - Threat from travelers with laptops, cell phones, ...

#### **5 Minute Break**

#### **Questions Before We Proceed?**

## **Getting Around Firewalls**
# **Subverting Firewalls**

- Along with possible bugs, packet filters have a fundamentally limited semantic model
  - They lack a full understanding of the meaning of the traffic they carry
    - In part because operate only at layers 3 & 4; not 7
- How can a local user who wants to get around their site's firewall exploit this?
  - (Note: we're not talking about how an external attacker can escape a firewall's restrictions)
- One method of subversion: abuse ports
  - Who says that e.g. port 53/udp = DNS?
    - Why couldn't it be say Skype or BitTorrent?
    - Just requires that client & server agree on application protocol





































### Packet Sent to Remote Relay



### Packet Sent by Remote Relay











# **Hiding on Other Ports**

- Method #1: use port allocated to another service (how can this be detected?)
- Method #2: tunneling
  - Encapsulate one protocol inside another
  - Receiver of "outer" protocol *decapsulates* interior tunneled protocol to recover it
  - Pretty much **any** protocol can be tunneled over another (with enough effort)
- E.g., tunneling IP over SMTP (email)
  - Just need a way to code an IP datagram as an email message (either mail body or just headers)

### **Example: Tunneling IP over Email**

From: halo-nut@yoyodyne.com To: my-buddy@tunnel-decapsulators.R.us Subject: Here's my IP datagram IP-header-version: 4 IP-header-len: 5 IP-proto: 17 (UDP) IP-src: 7.7.8.2 IP-dst: 14.6.1.7 IP-payload: 0xa144bf2c0102...

Remote email server receives this legal email, **builds** an IP packet corresponding to description in email body ...

... and **injects** it into the network How can a firewall detect this??

### **Network Control & Tunneling**

- *Tunneling* = embedding one protocol inside another
  - Sender and receiver at each side of the tunnel both cooperate (so it's not useful for initial attacks)
- Traffic takes on properties of outer protocol
  - Including for firewall inspection, which generally can't analyze inner protocol (due to complexity)
- Tunneling has legitimate uses
  - E.g., Virtual Private Networks (VPNs)
    - Tunnel server relays remote client's packets
    - Makes remote machine look like it's local to its home network
    - Tunnel encrypts traffic for privacy & to prevent meddling

### Other Ways of Securing Network Access

#### **Secure External Access to Inside Machines**



- Often need to provide secure remote access to a network protected by a firewall
  - Remote access, telecommuting, branch offices, ...
- Create secure channel (*Virtual Private Network*, or VPN) to tunnel traffic from outside host/network to inside network
  - Provides Authentication, Confidentiality, Integrity
    - Requires some form of key management to set up
  - However, also raises *perimeter issues* 
    - (Try it yourself at http://www.net.berkeley.edu/vpn/)

## **Application Proxies**

- Can more directly control applications by requiring them to go through a proxy for external access
  - Proxy doesn't simply forward, but acts as an application-level middleman
- Example: SSH gateway
  - Require all SSH in/out of site to go through gateway
  - Gateway logs authentication, inspects decrypted text
  - Site's firewall configured to prohibit any other SSH access

### **SSH Gateway Example**



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- Example: SSH gateway
  - Require all SSH in/out of site to go through gateway
  - Gateway logs authentication, inspects decrypted text
  - Site's firewall configured to prohibit any other SSH access
- Provides a powerful degree of monitoring/control
- Costs?
  - Need to run extra server(s) per app (possible *bottleneck*)
  - Each server requires careful hardening