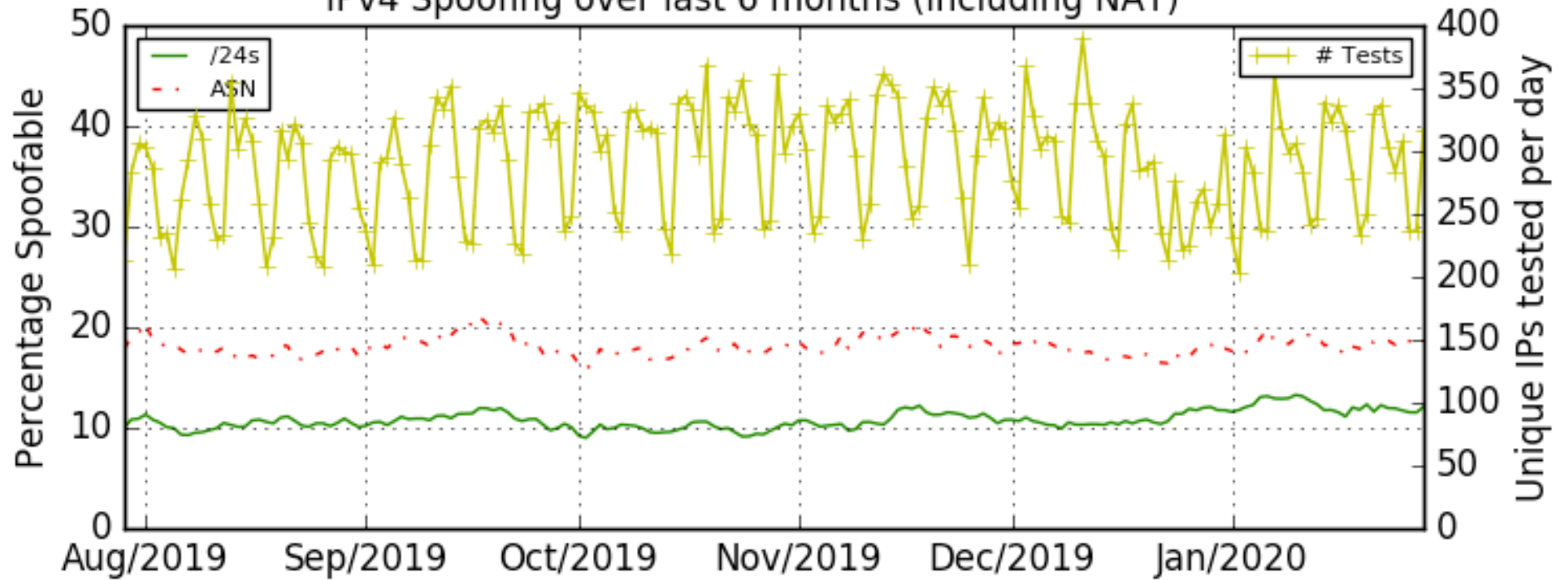
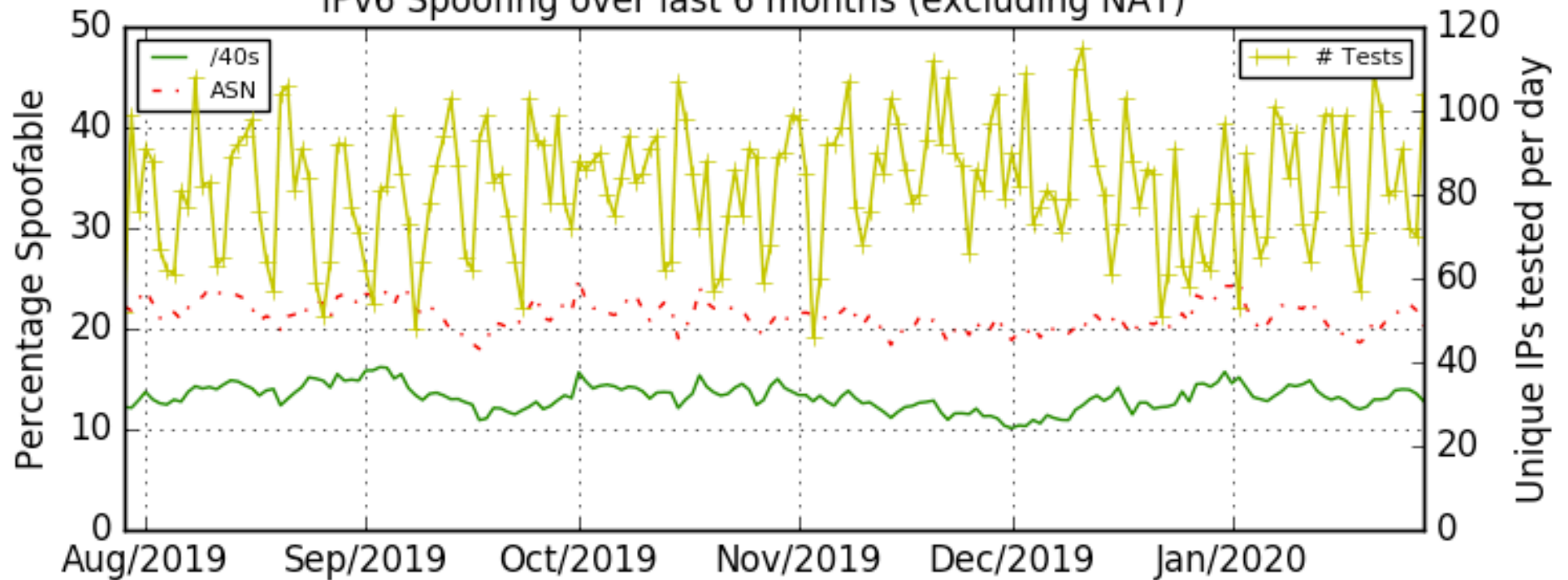
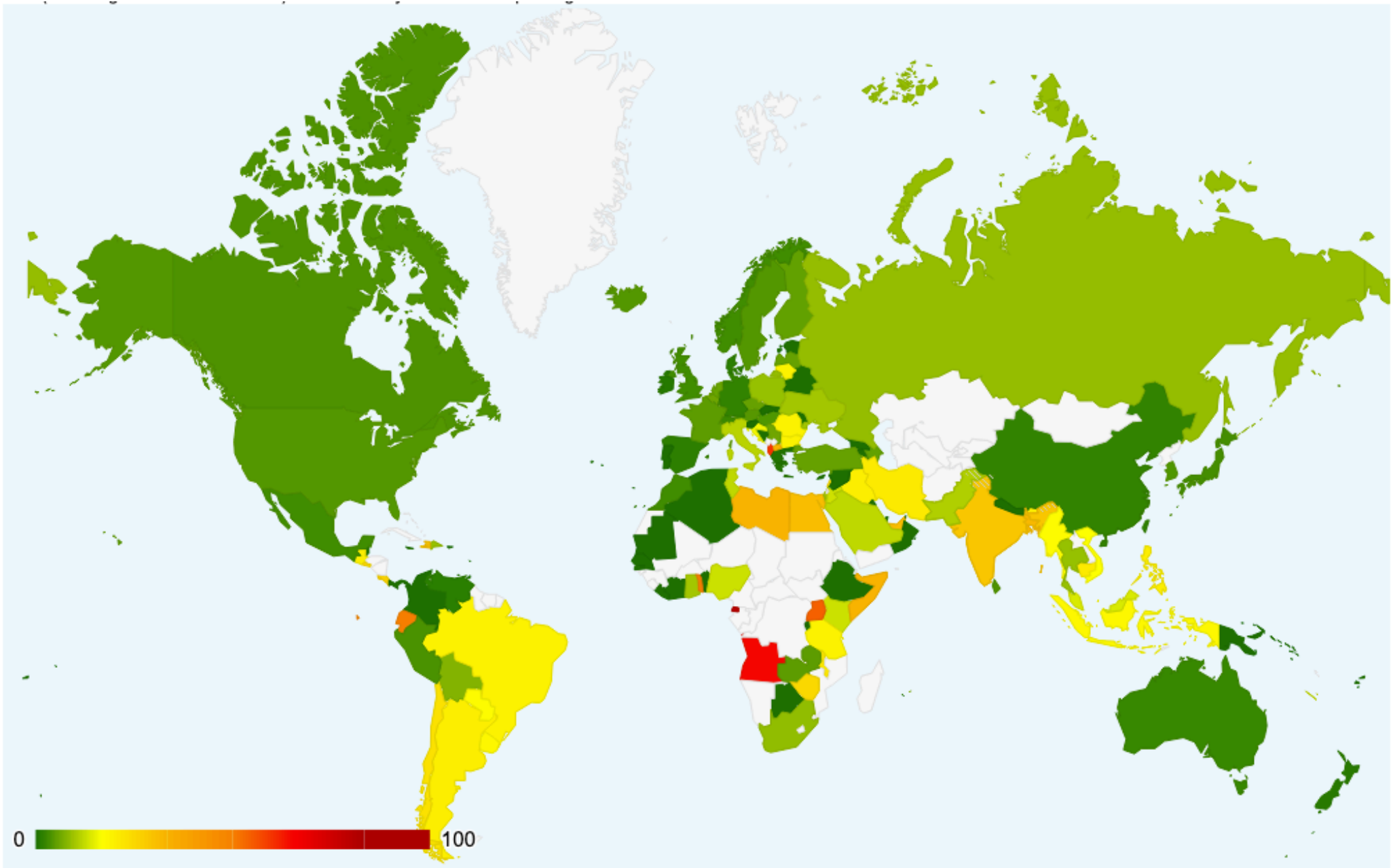


IPv4 Spoofing over last 6 months (including NAT)

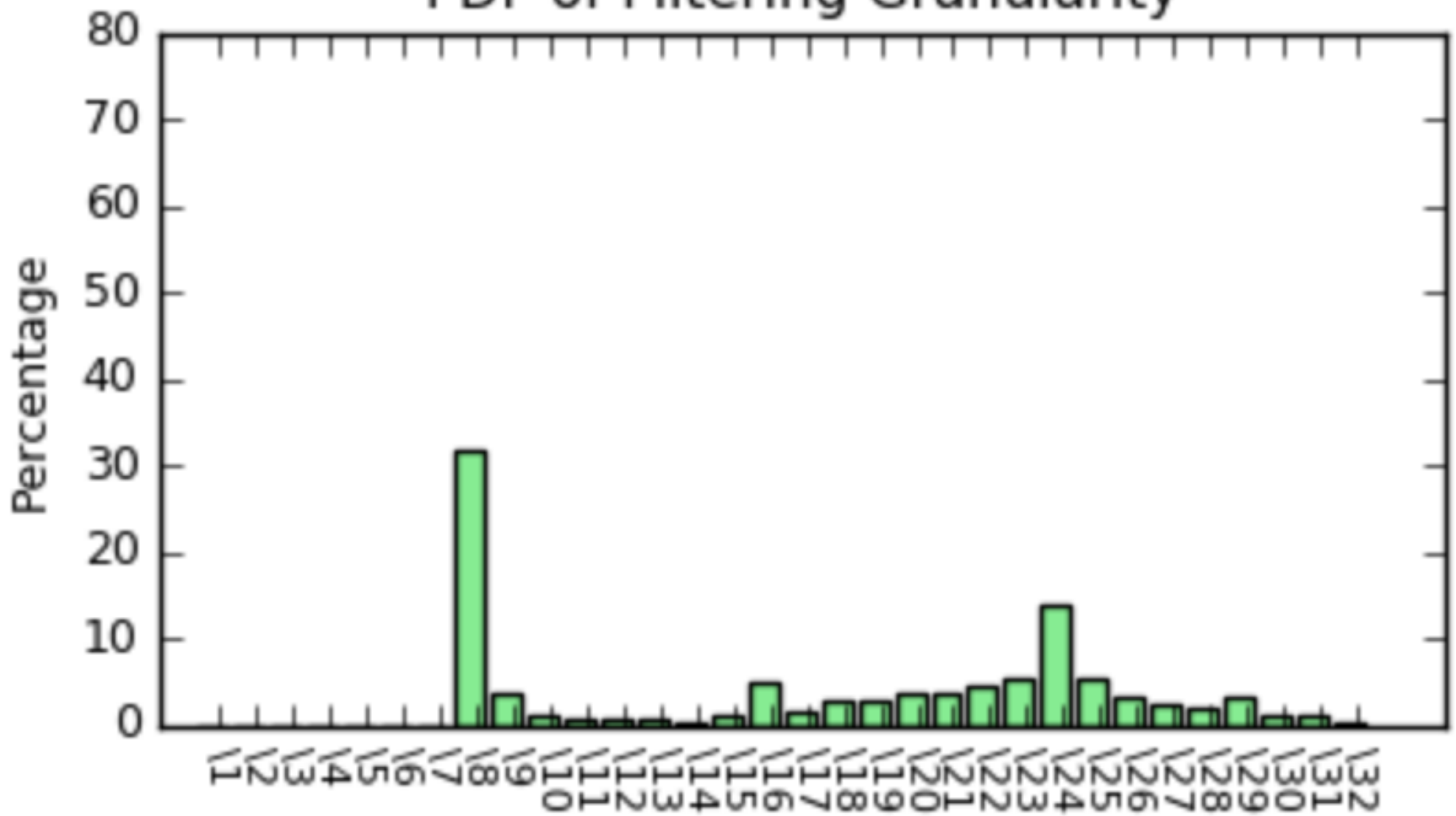


IPv6 Spoofing over last 6 months (excluding NAT)

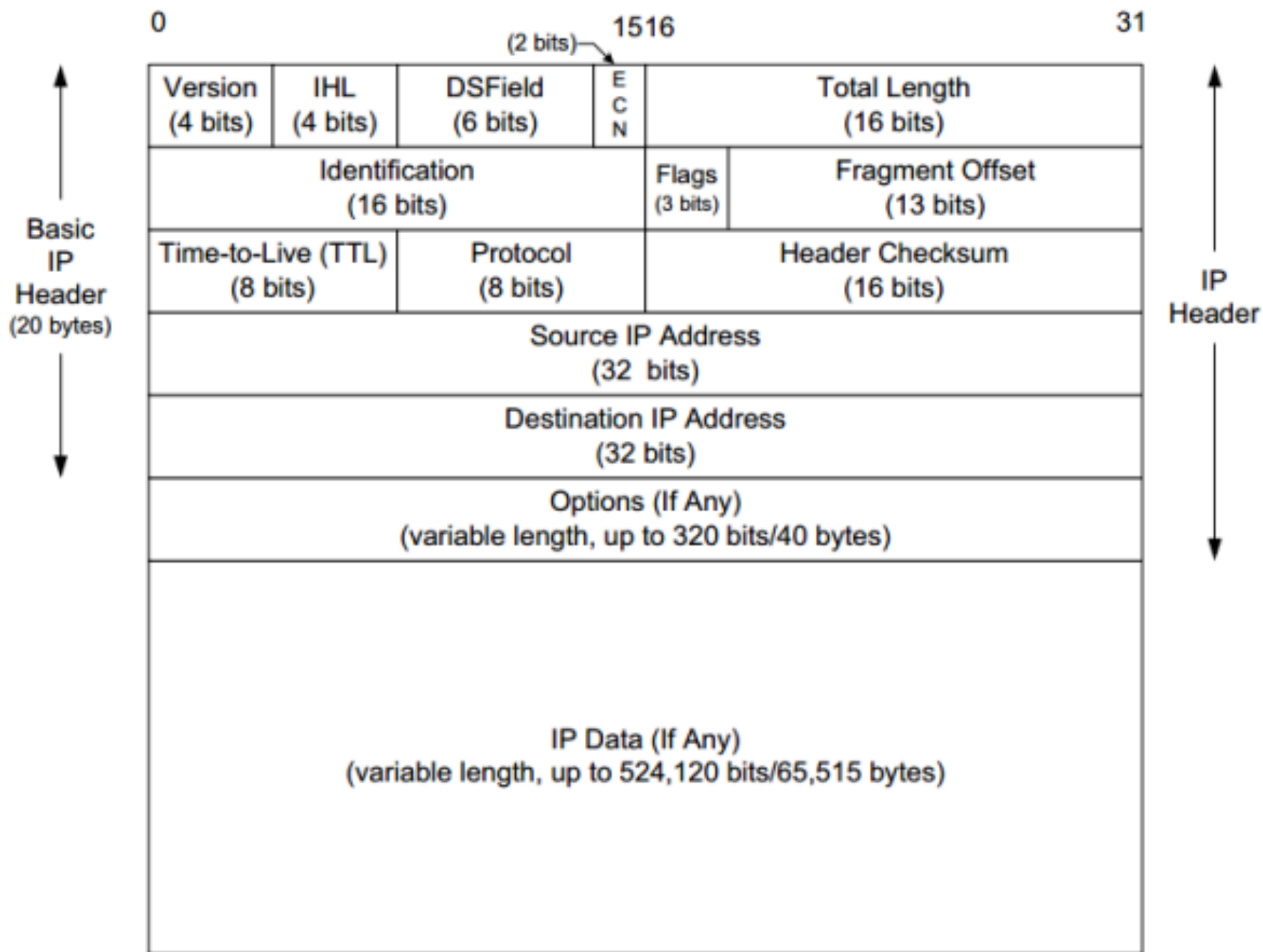


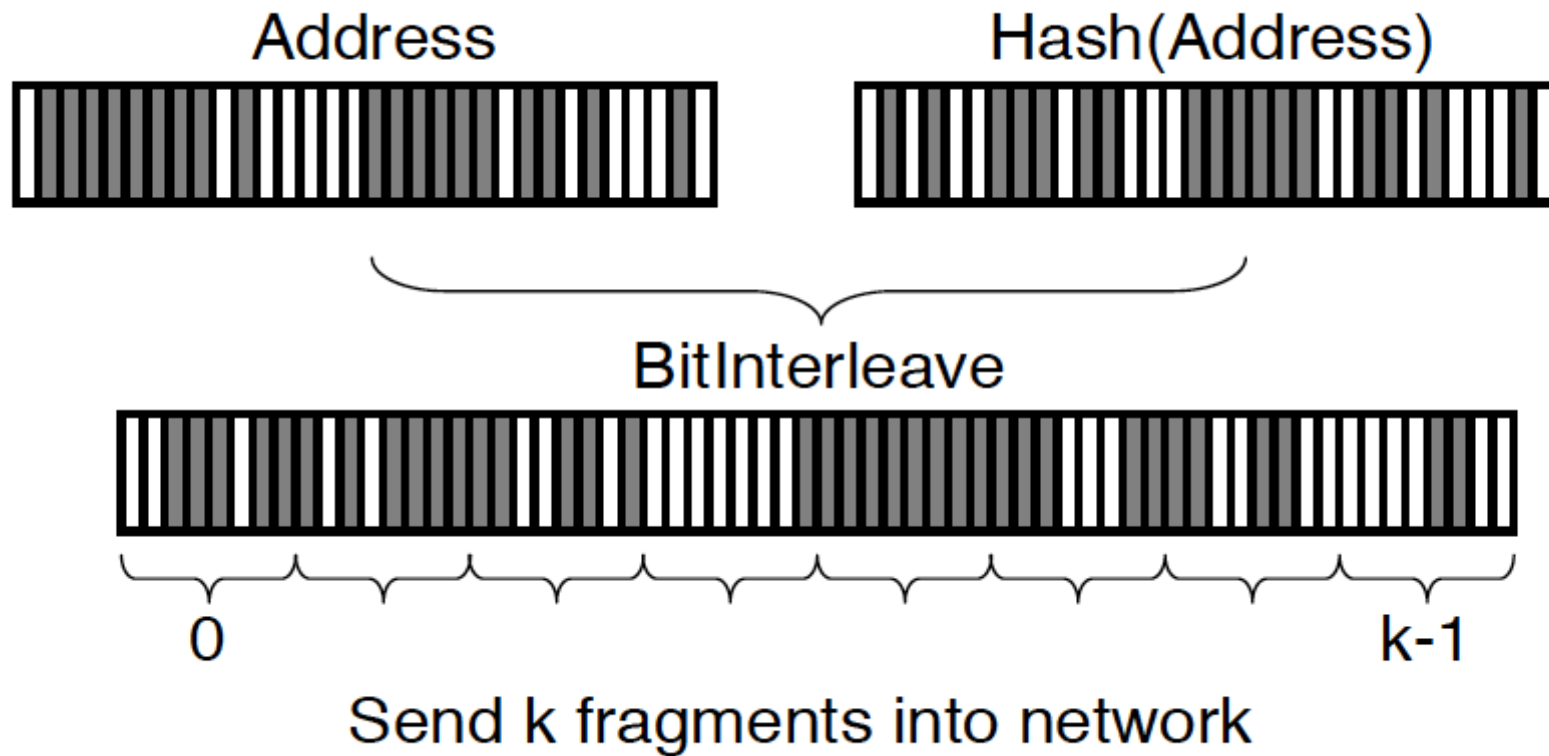


PDF of Filtering Granularity

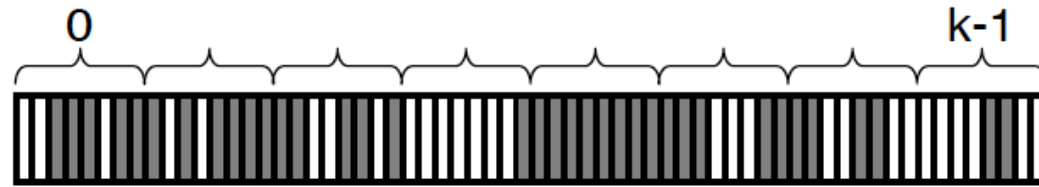


We define the *approximate trace-back* problem as finding a candidate attack path for each attacker that contains the true attack path as a suffix. We call this the *valid suffix* of the candidate path.





Combine k fragments from network



BitDeinterleave



Address?

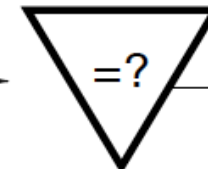


Hash(Address)?

Hash



Hash(Address?)



No

Reject

Yes



Address

for $d := 0$ to $maxd$

for all ordered combinations of fragments at distance d

construct edge z

if $d \neq 0$ then

$z := z \oplus last$

if $Hash(EvenBits(z)) = OddBits(z)$ then

insert edge $(z, EvenBits(z), d)$ into G

$last := EvenBits(z);$

	Management overhead	Network overhead	Router overhead	Distributed capability	Post-mortem capability	Preventative/ reactive
Ingress filtering	Moderate	Low	Moderate	N/A	N/A	Preventative
Link testing						
Input debugging	High	Low	High	Good	Poor	Reactive
Controlled flooding	Low	High	Low	Poor	Poor	Reactive
Logging	High	Low	High	Excellent	Excellent	Reactive
ICMP Traceback	Low	Low	Low	Good	Excellent	Reactive
Marking	Low	Low	Low	Good	Excellent	Reactive

Table 1: Qualitative comparison of existing schemes for combating anonymous attacks and the probabilistic marking approach we propose.